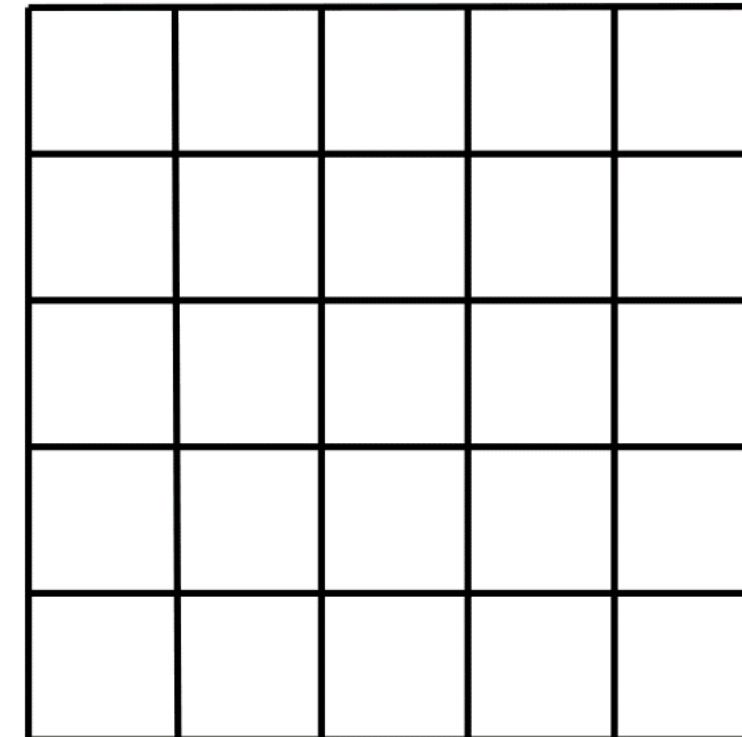
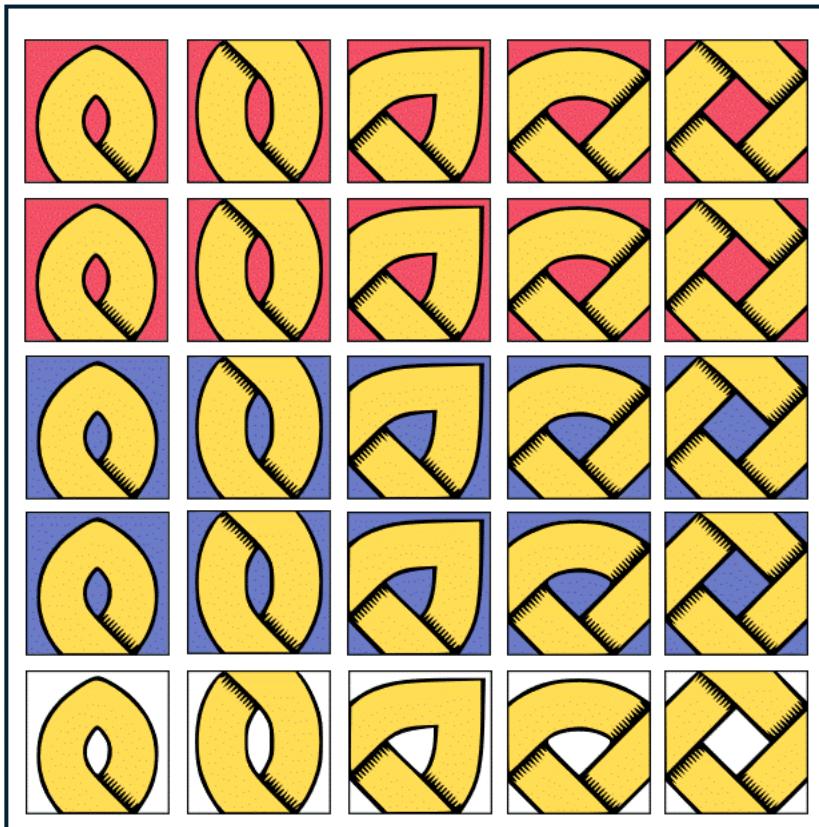
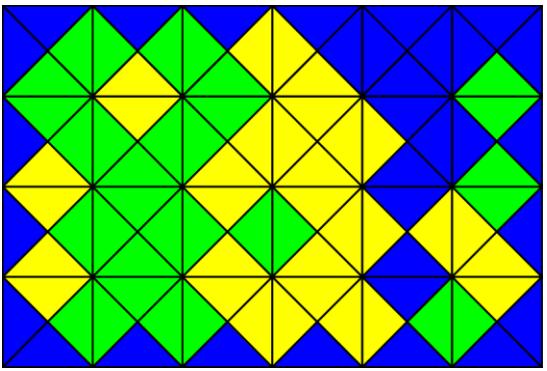


Tile-based Knot Assembly with Celtic!

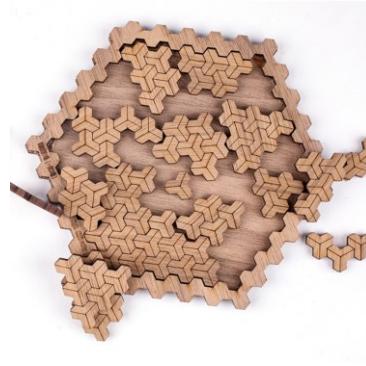


Authors: Divya Bajaj, **Ryan Knobel**, Juan Manuel Perez, Rene Reyes,
Ramiro Santos, Tim Wylie

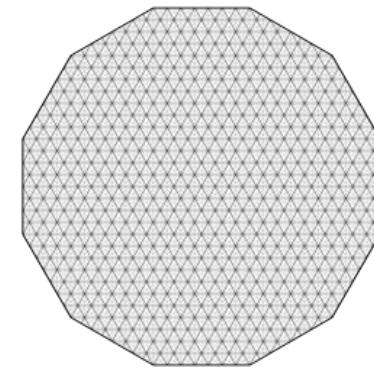
Background



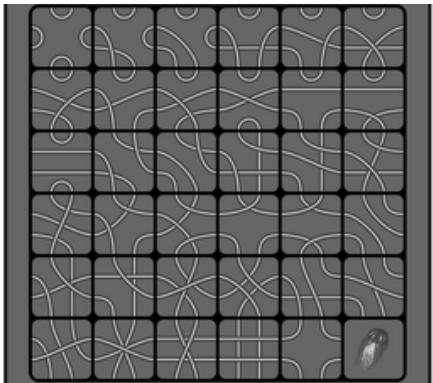
MacMahon Squares



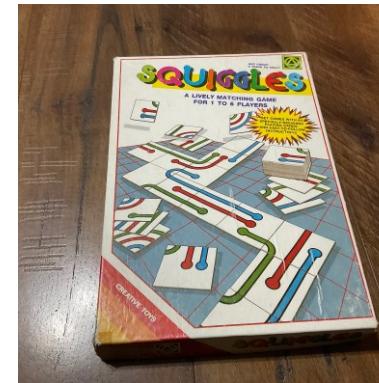
Jigsaw Puzzles



Eternity

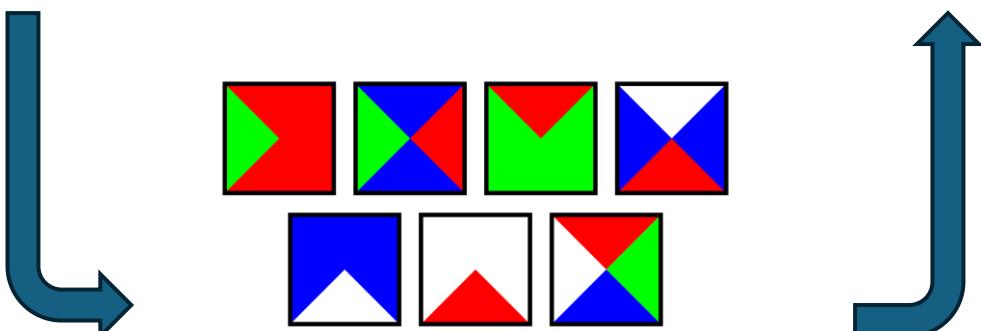
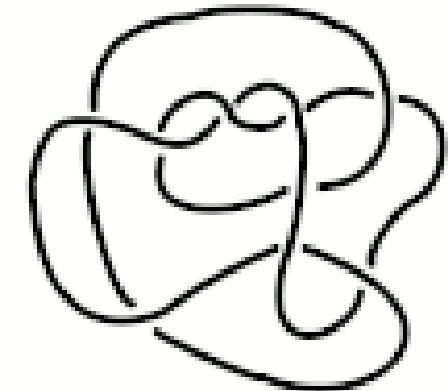
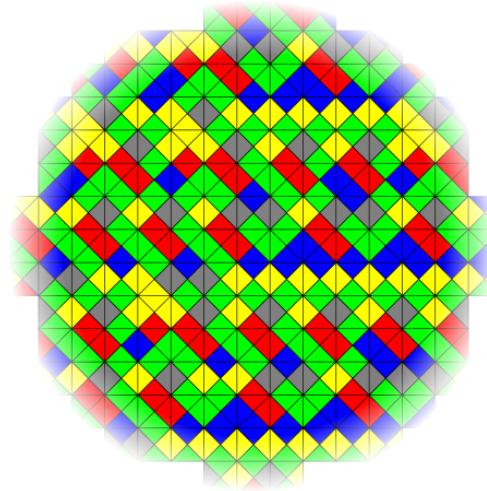
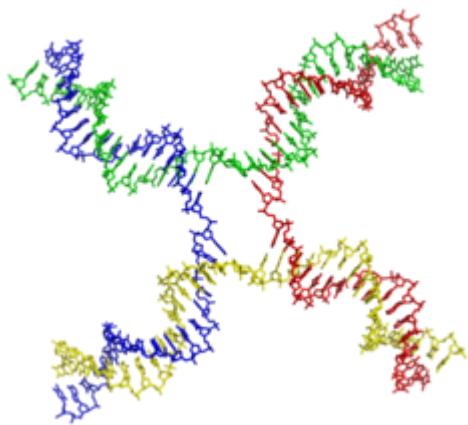


Tsuro



Squiggles

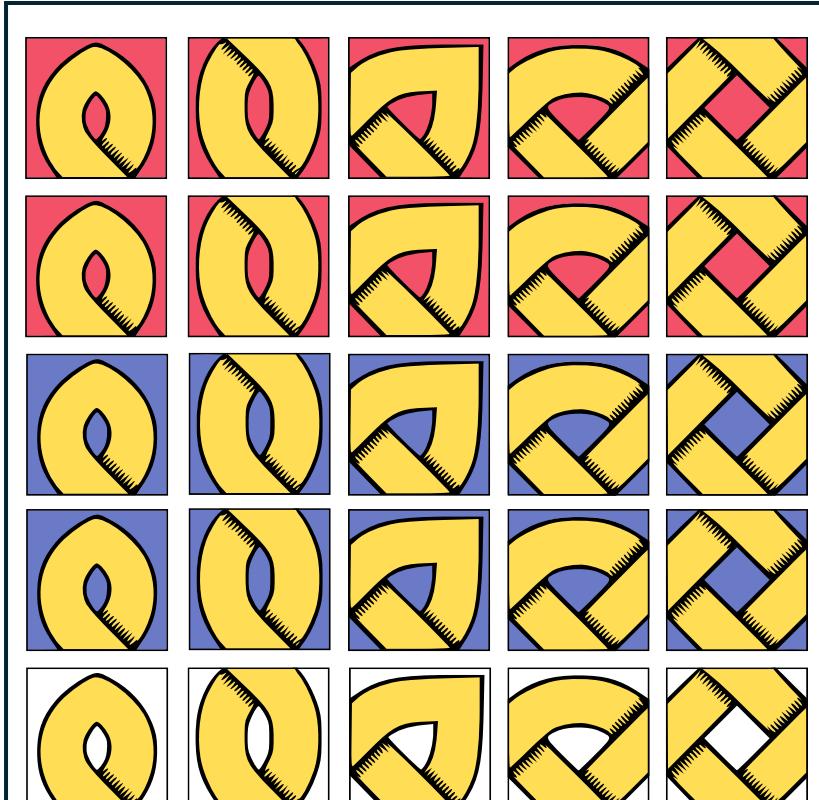
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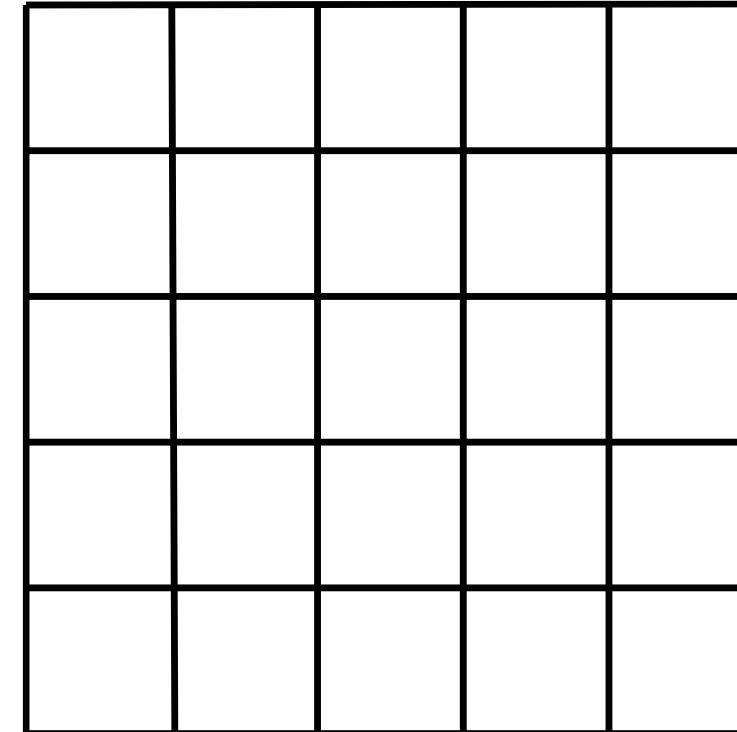
Wang tiles

[Knot Theory Wikipedia]

Celtic!

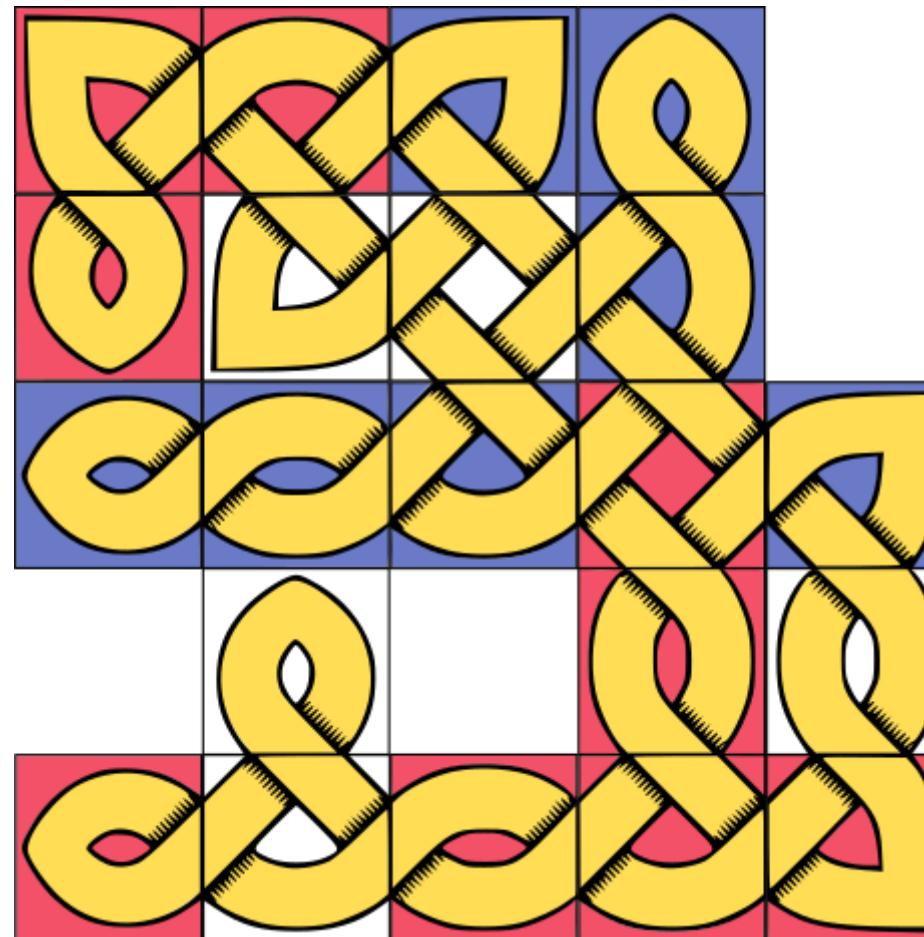


Pieces

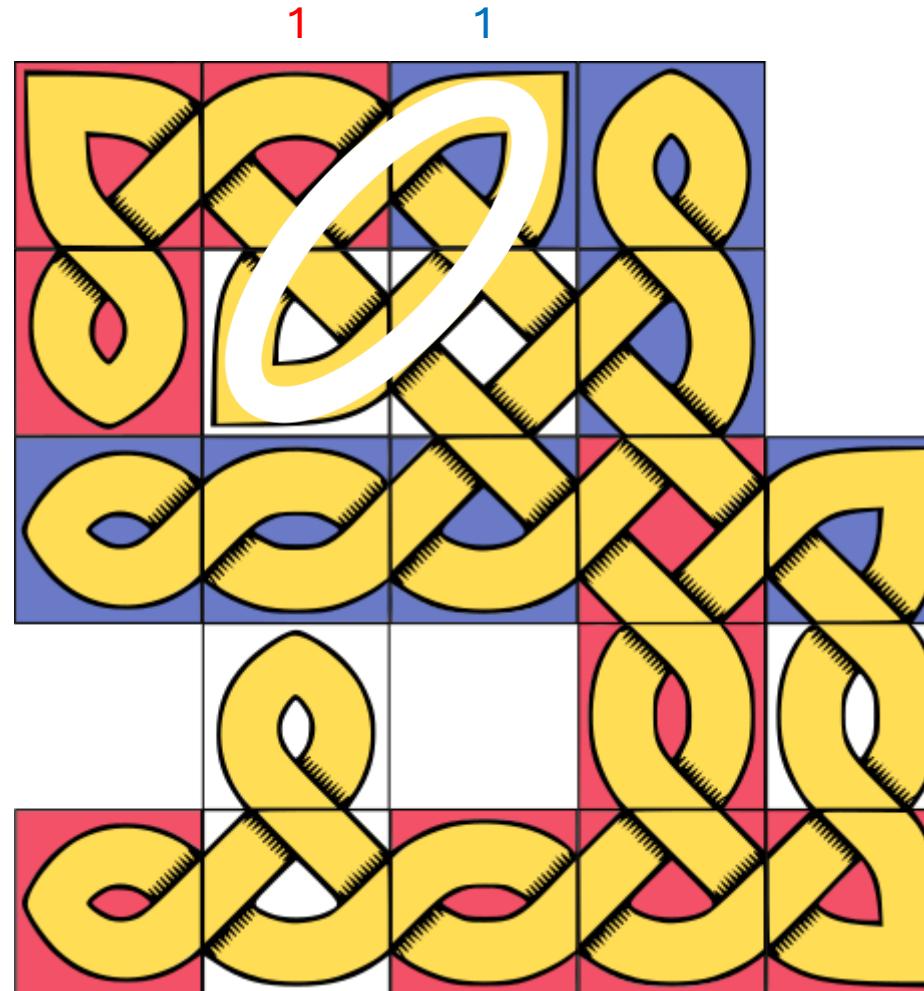


Board

Celtic! - Scoring



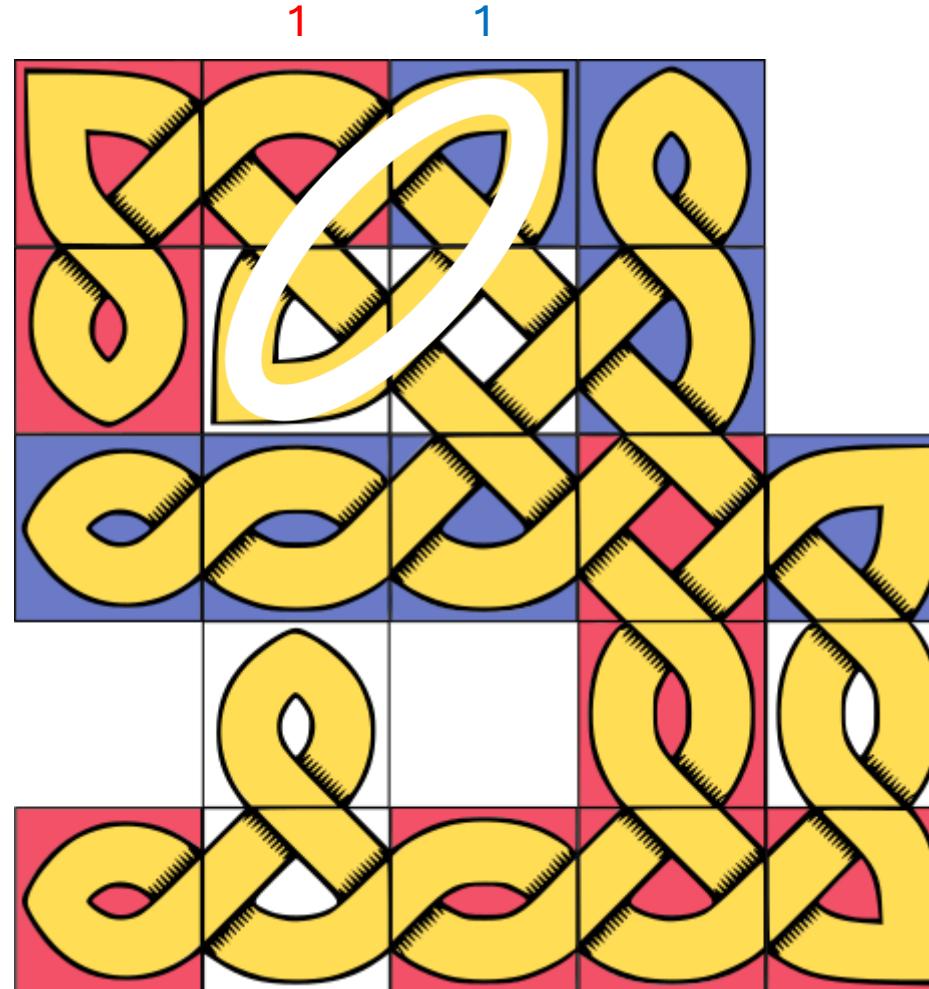
Celtic! - Scoring



Celtic! - Scoring

Score

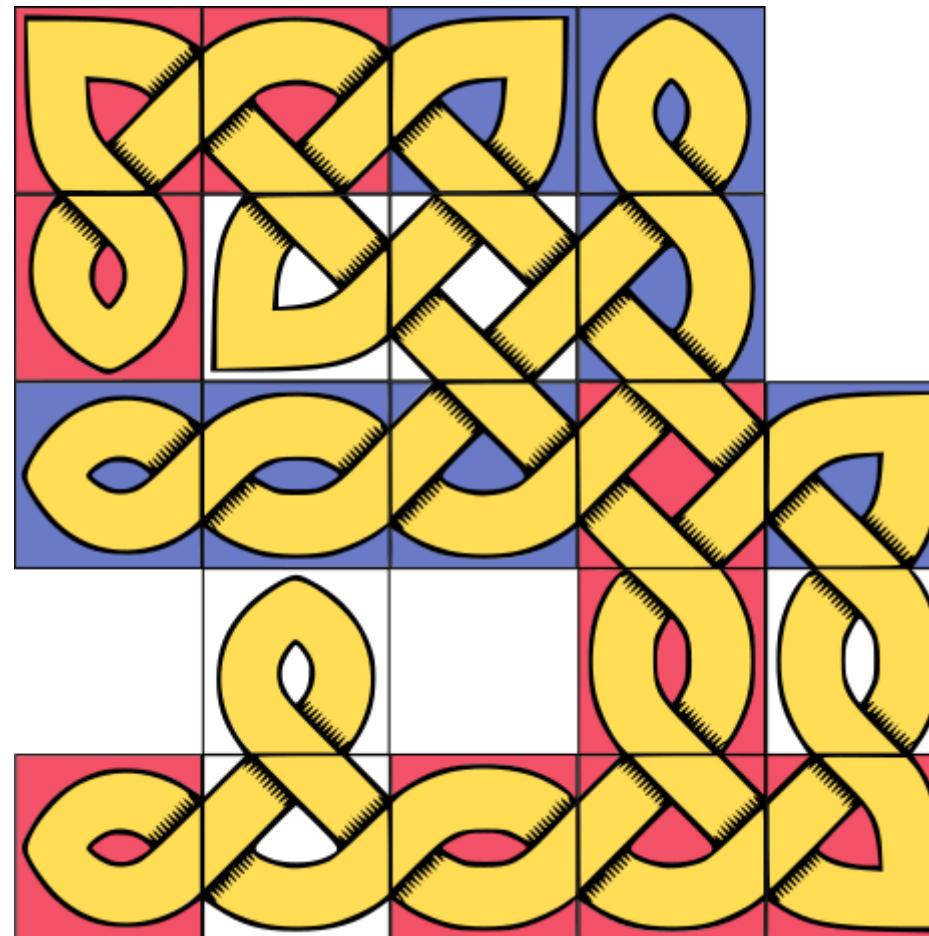
Knot 1: Tie



Celtic! - Scoring

Score

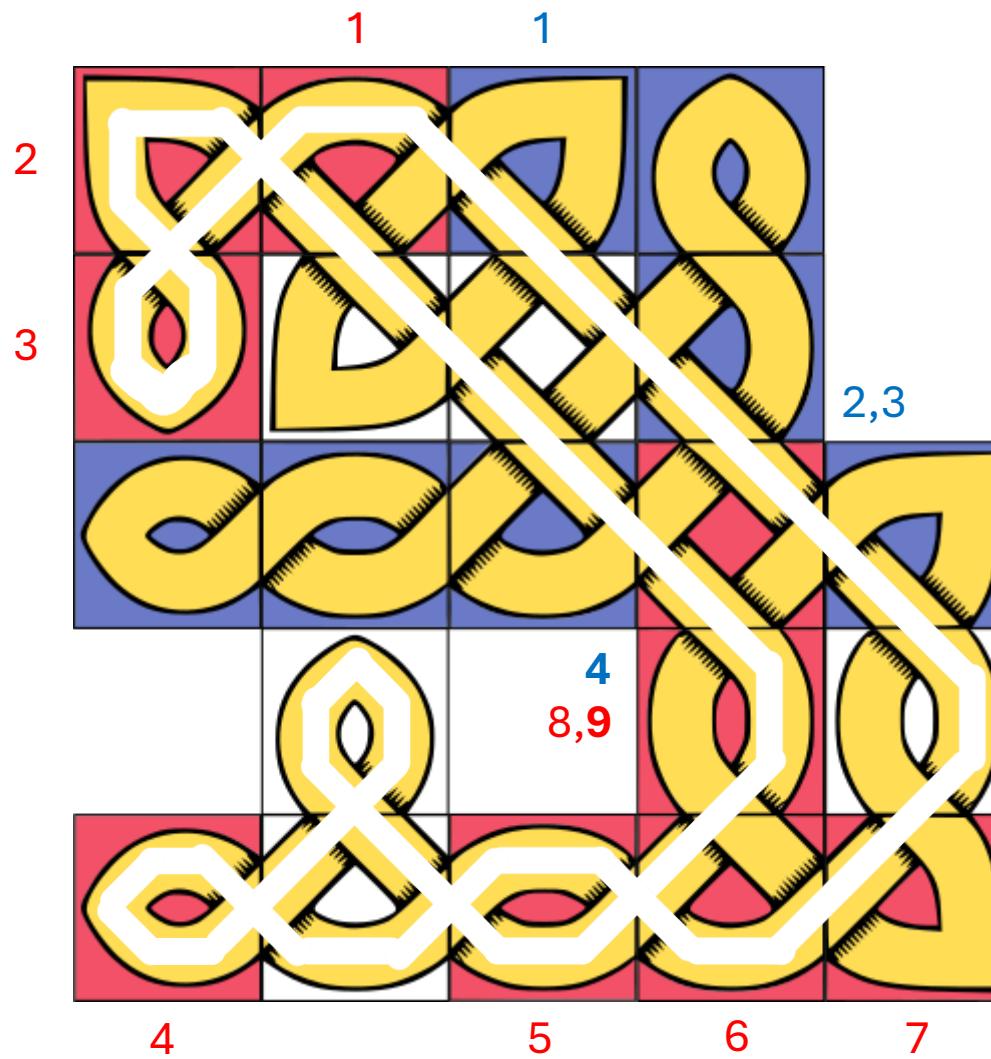
Knot 1: Tie



Celtic! - Scoring

Score

Knot 1: Tie

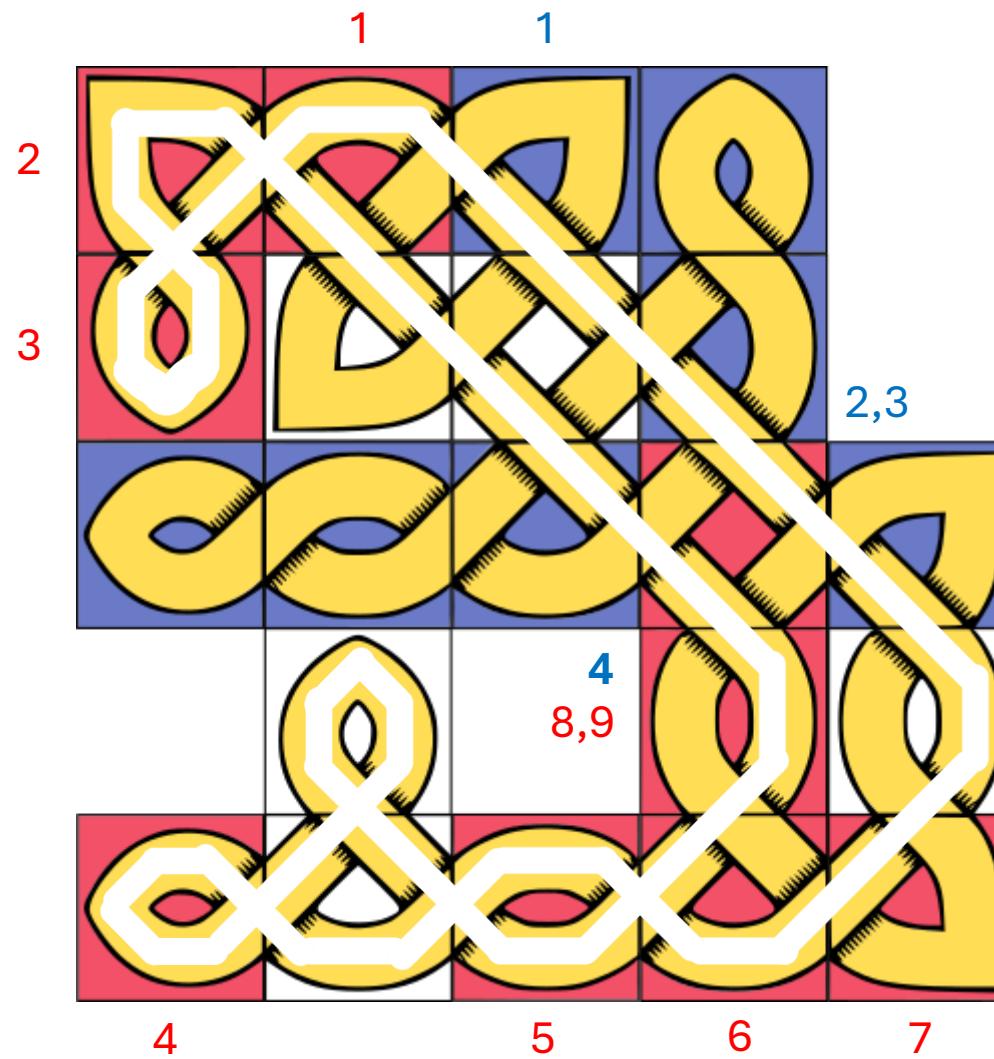


Celtic! - Scoring

Score

Knot 1: Tie

Knot 2: Red (9)

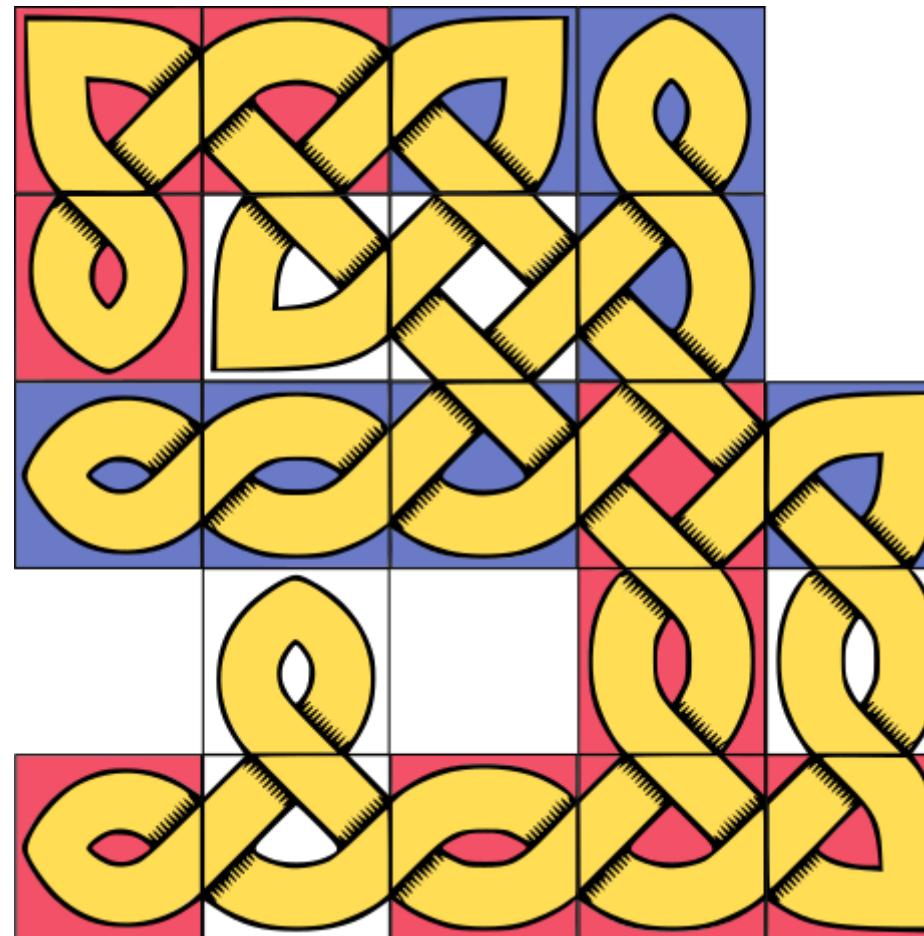


Celtic! - Scoring

Score

Knot 1: Tie

Knot 2: Red (9)



Celtic! - Scoring

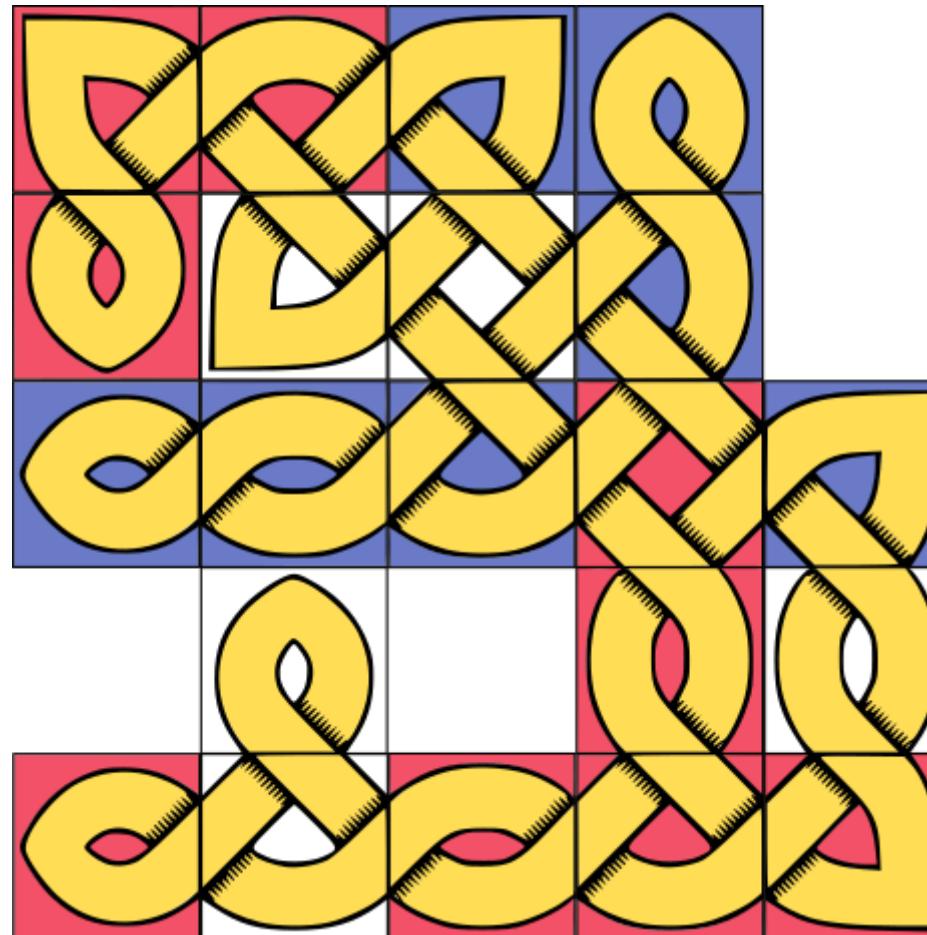
Score

Knot 1: Tie

Knot 2: Red (9)

Knot 3: Blue (5)

Knot 4: Red (4)



Celtic! - Scoring

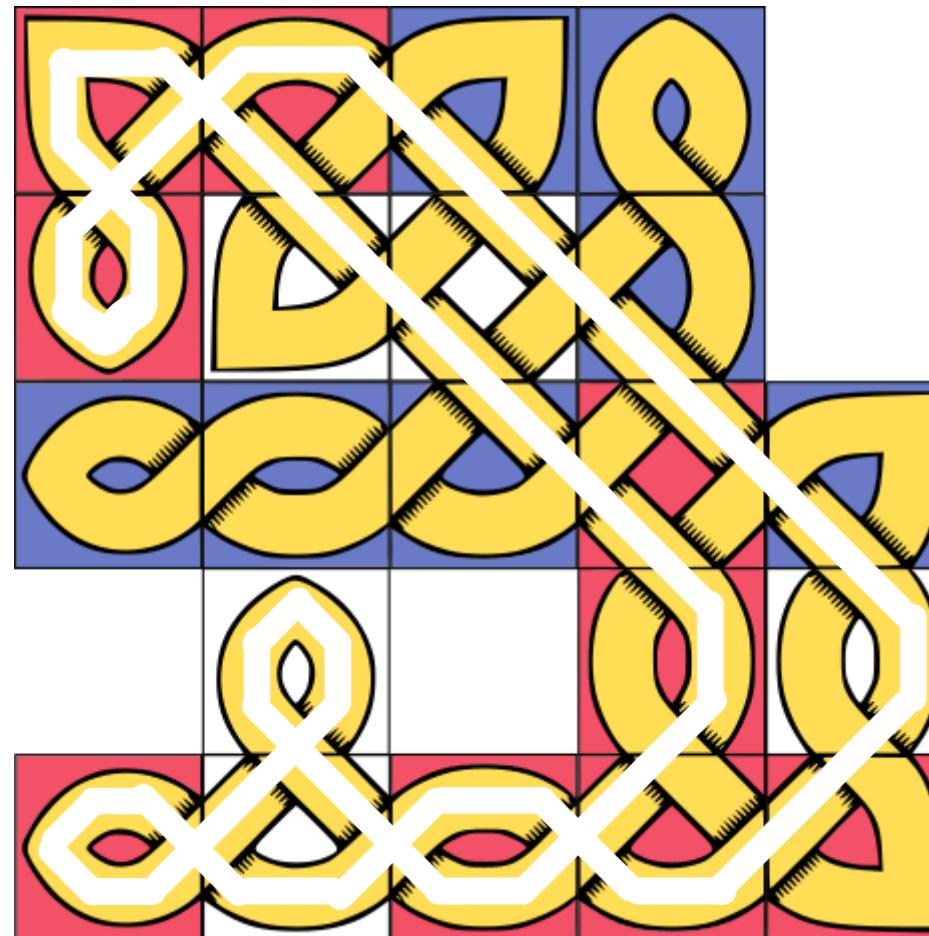
Score

Knot 1: Tie

Knot 2: Red (9)

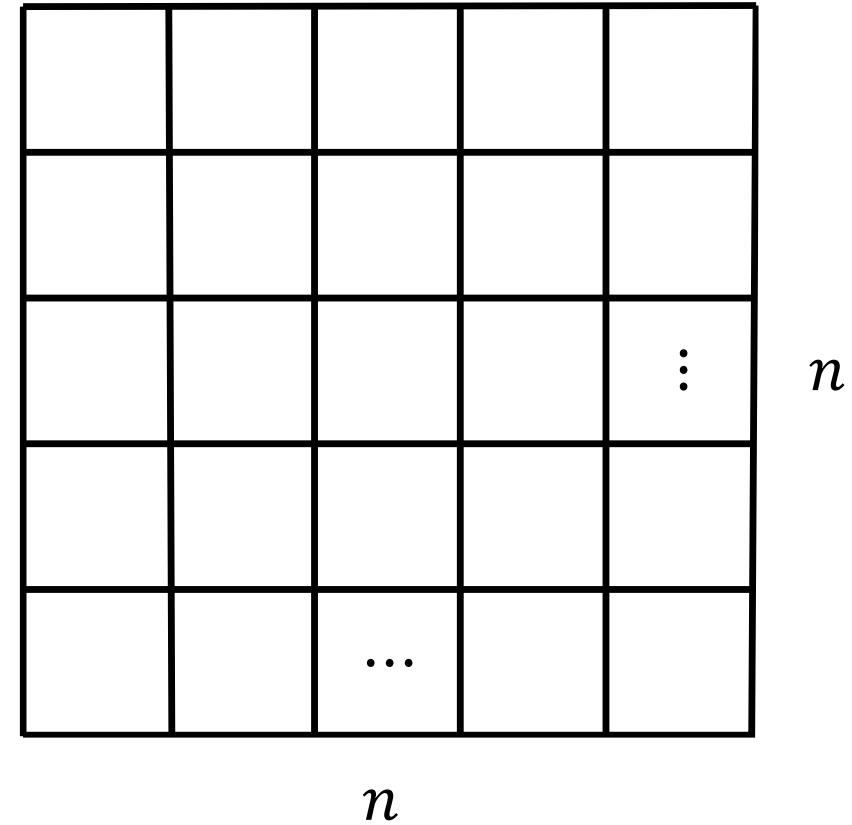
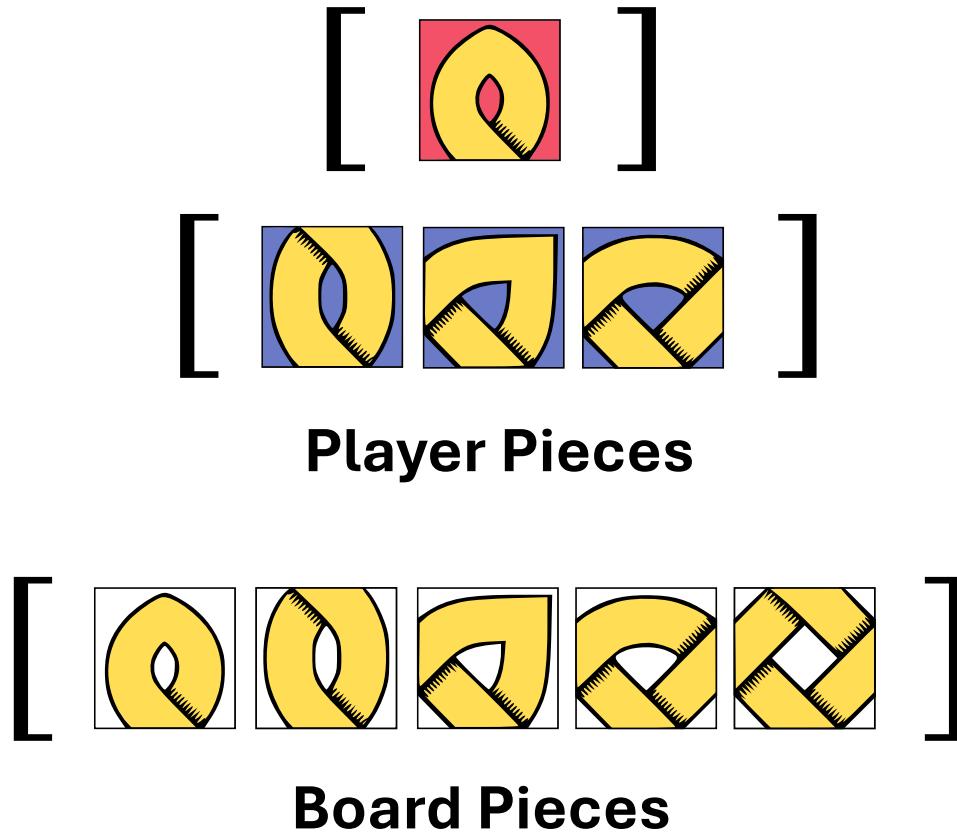
Knot 3: Blue (5)

Knot 4: Red (4)



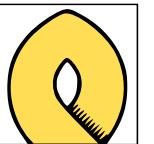
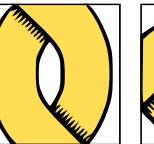
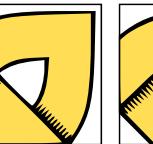
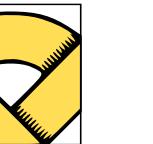
Generalized 2-Player Celtic!

Generalization: 2-Player



Problem: Does the blue player have a winning strategy?

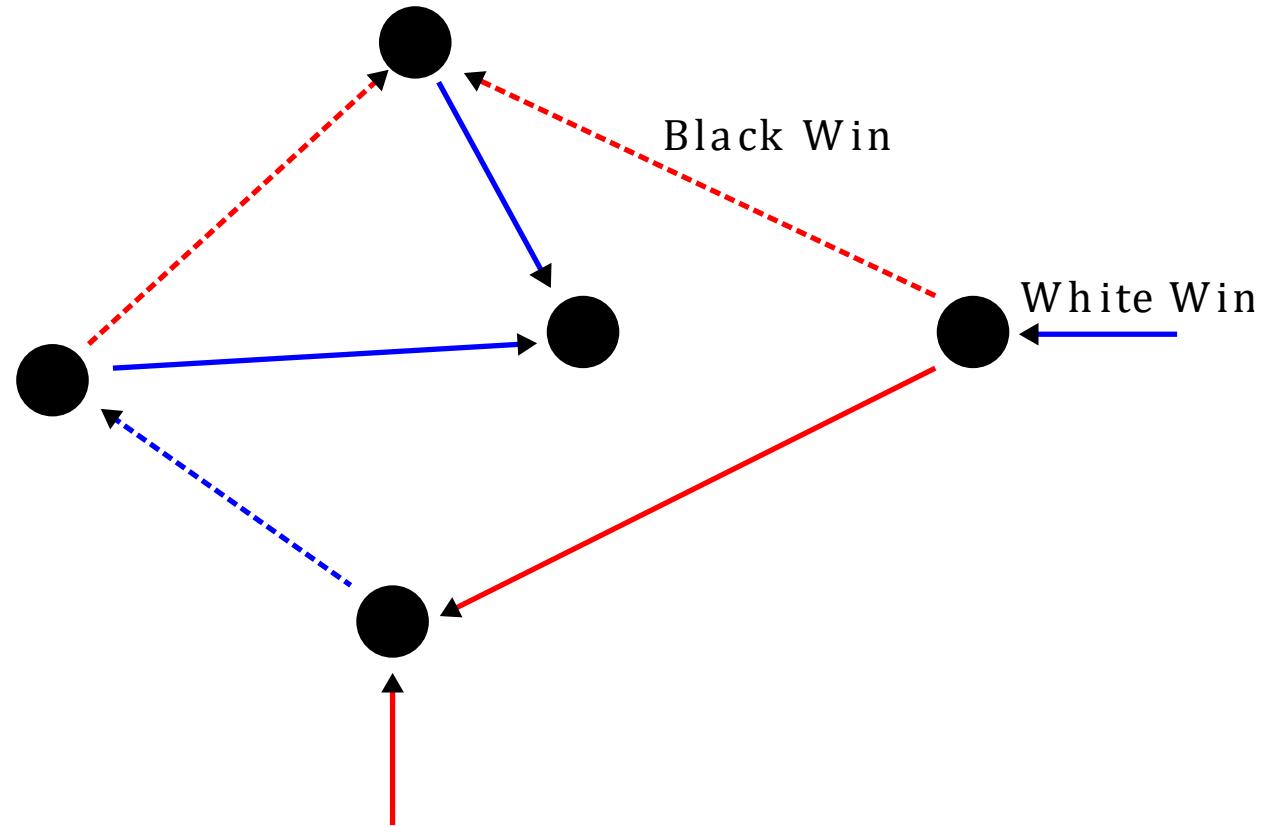
2-Player Celtic!

Player Pieces	Board Pieces	Complexity
   	   	<i>PSPACE-complete</i>

Problem: Does the blue player have a winning strategy?

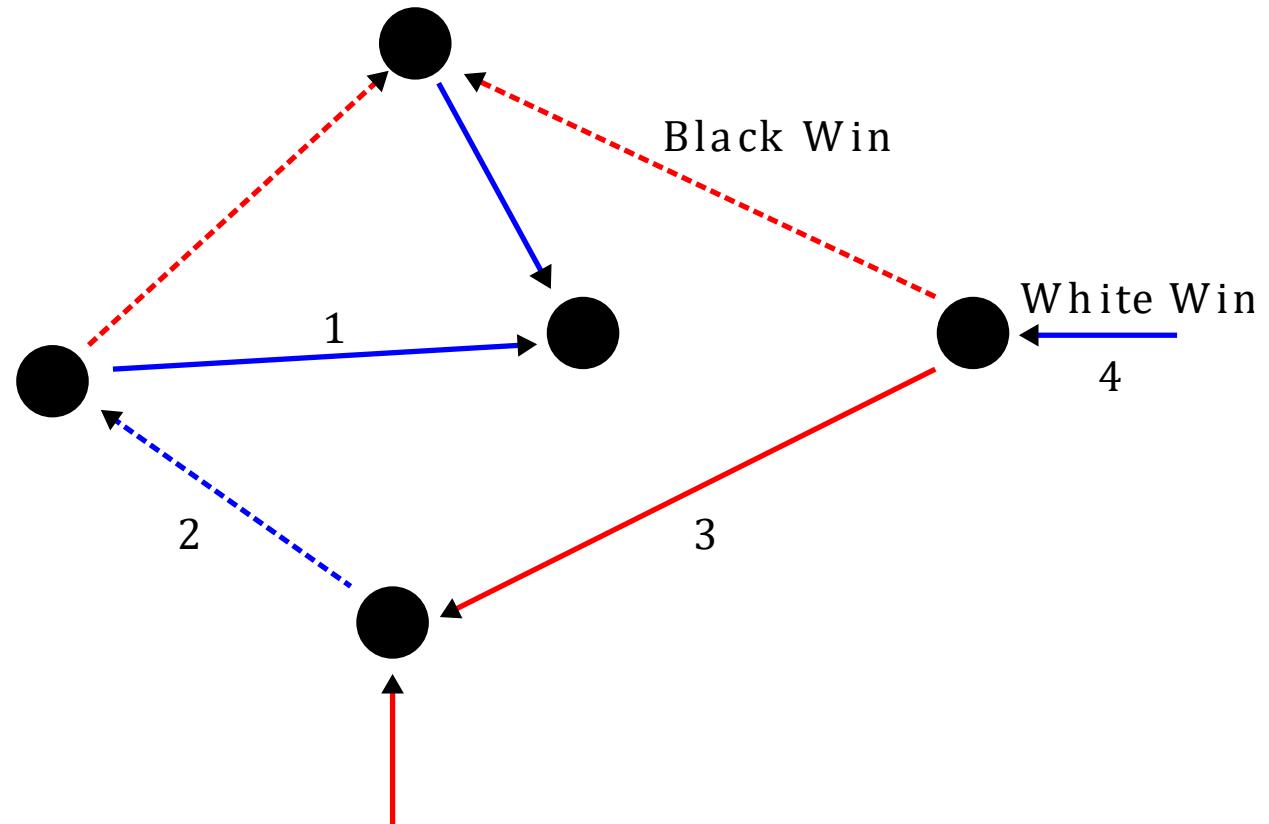
2-Player Celtic!

- Bounded 2-Player Constraint Logic
- Given directed graph with edges of weights 1 or 2
- Edges are partitioned for Black (dotted)/White (solid) players
- Goal is to flip a specific edge

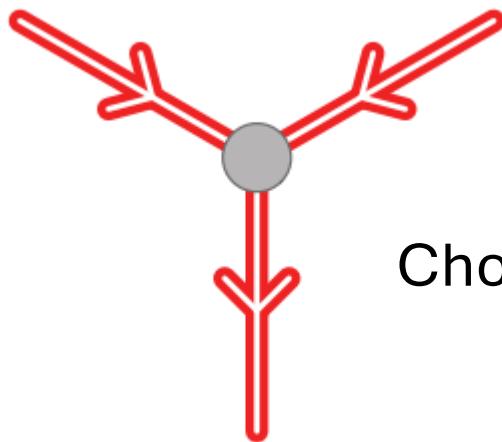


2-Player Celtic!

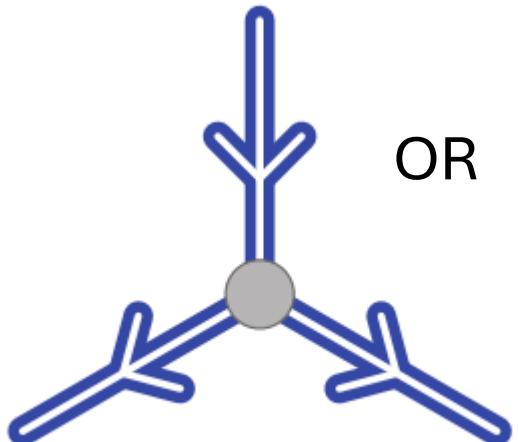
- Bounded 2-Player Constraint Logic
- Given directed graph with edges of weights 1 or 2
- Edges are partitioned for Black (dotted)/White (solid) players
- Goal is to flip a specific edge



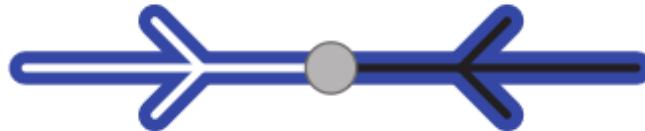
2-Player Celtic!



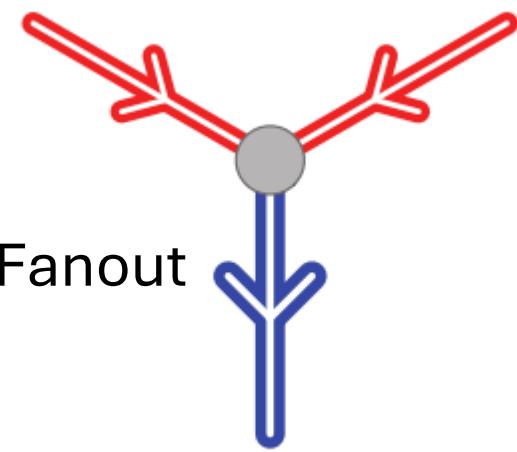
Choice



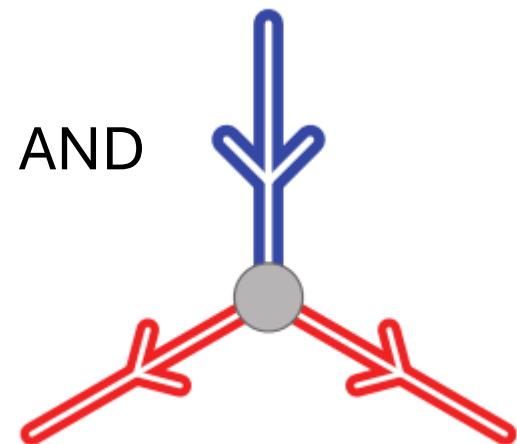
OR



Variable

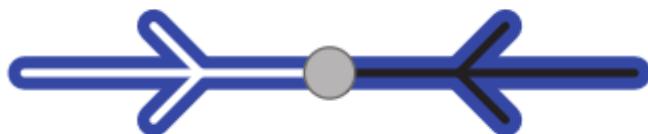
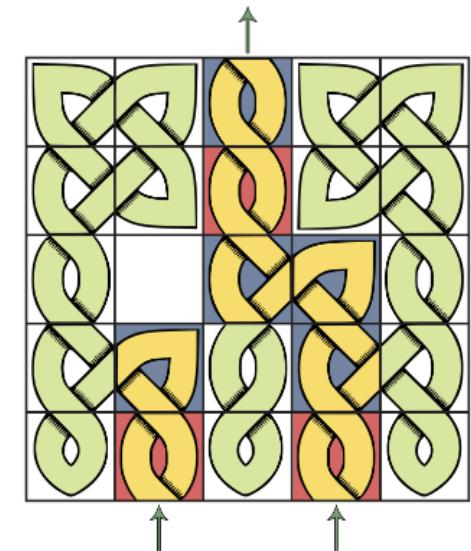
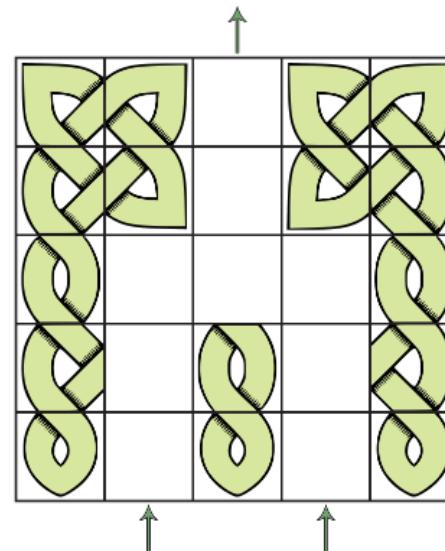
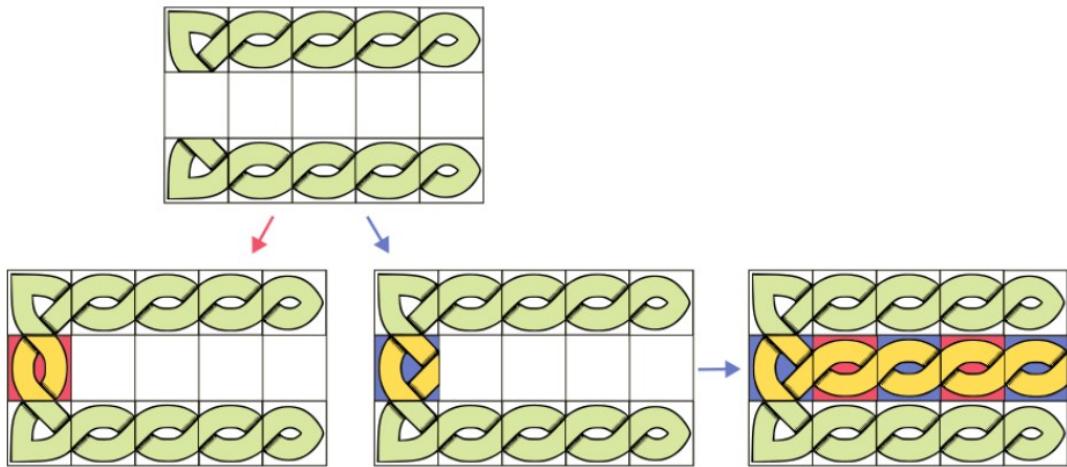


Fanout

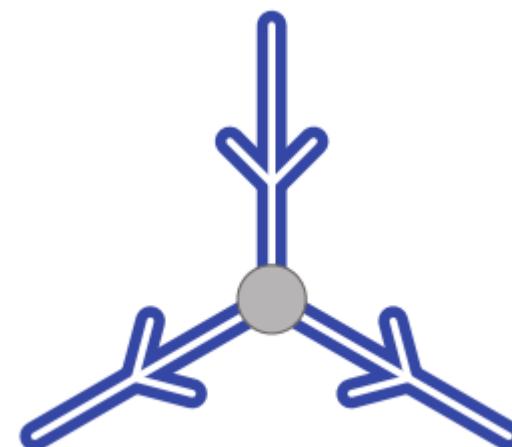


AND

2-Player Celtic!

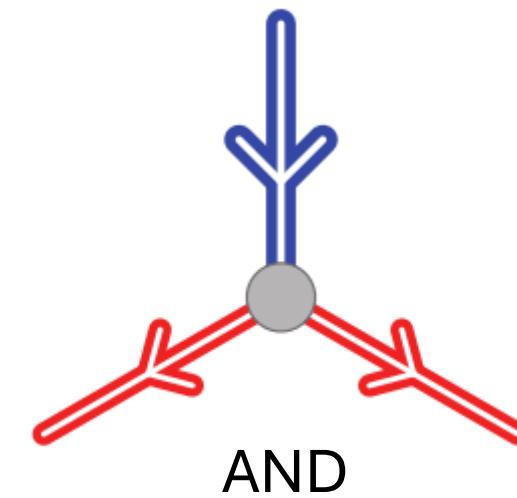
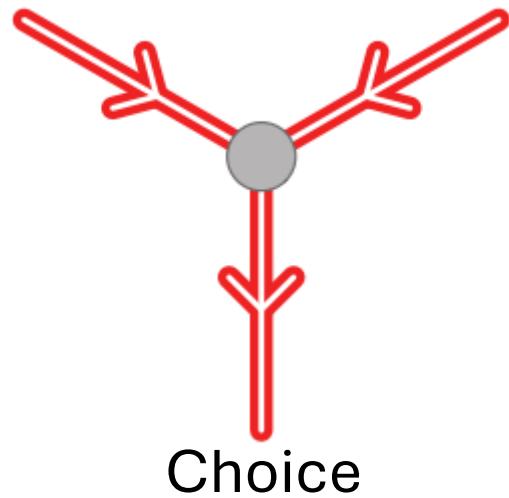
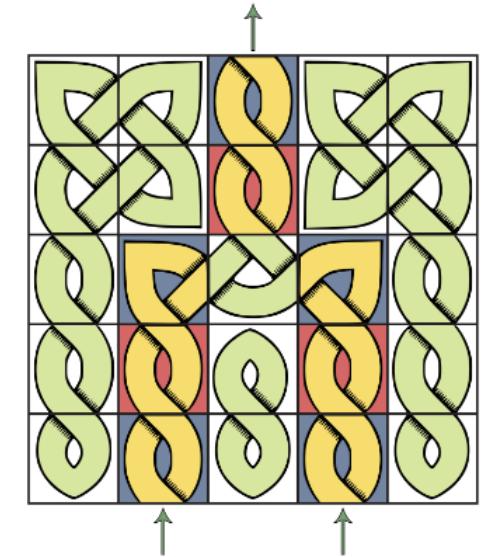
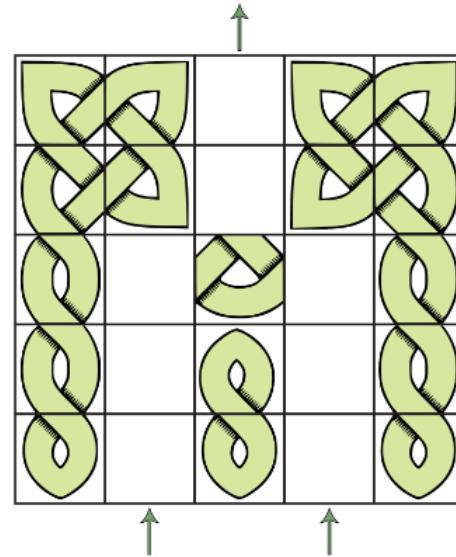
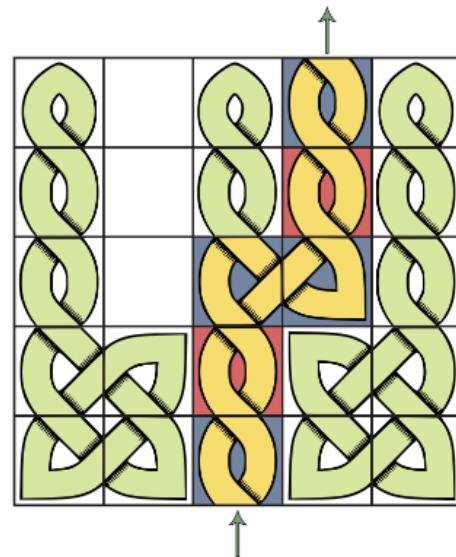
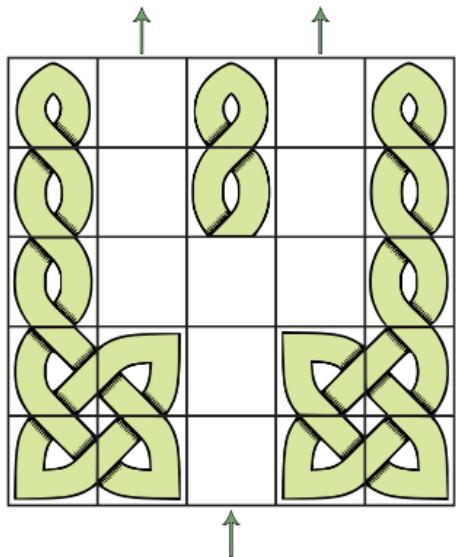


Variable

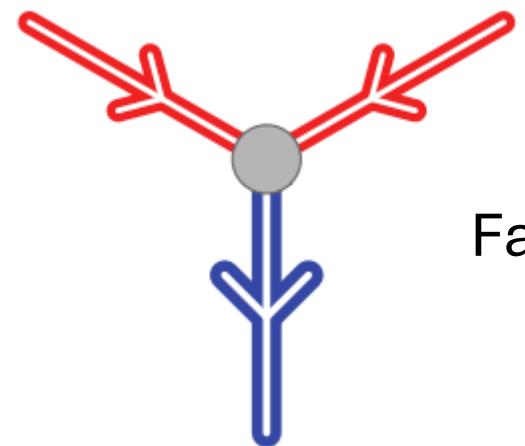
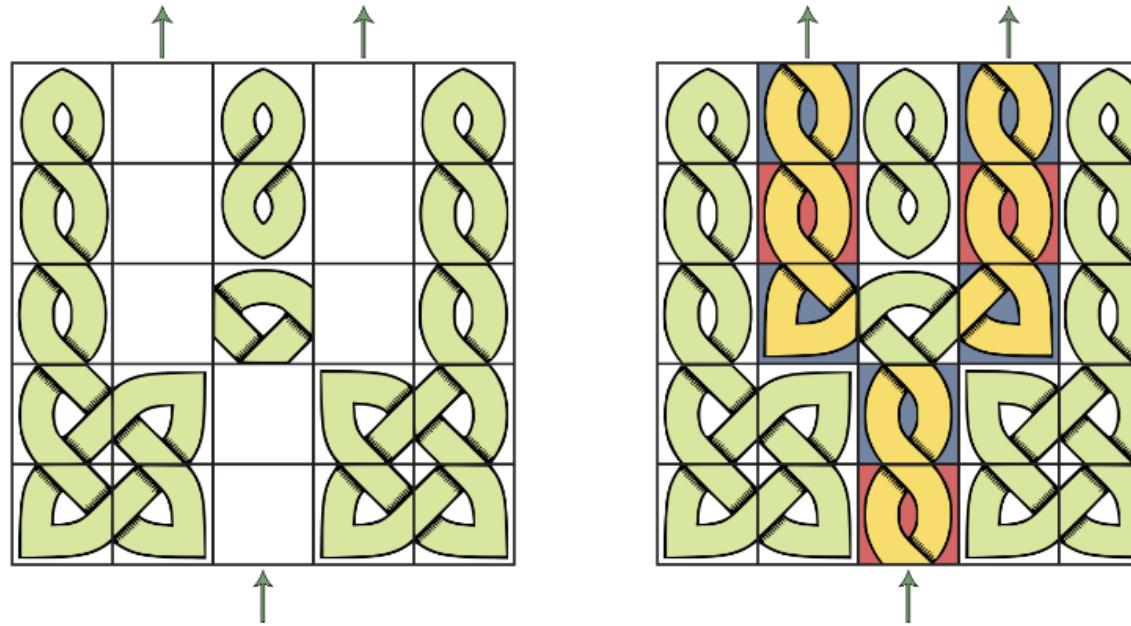


OR

2-Player Celtic!

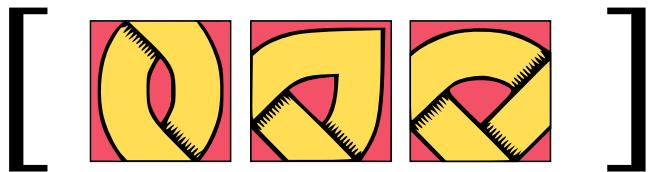


2-Player Celtic!

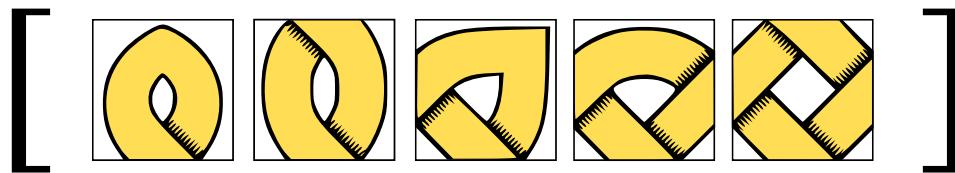


Generalized 1-Player Celtic!

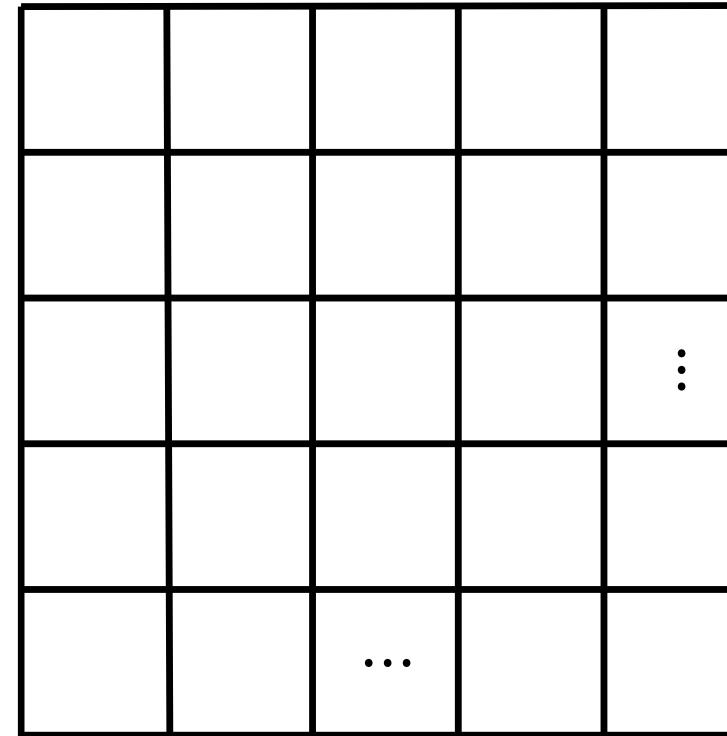
Generalization: 1-Player



Player Pieces



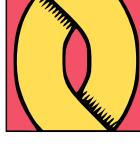
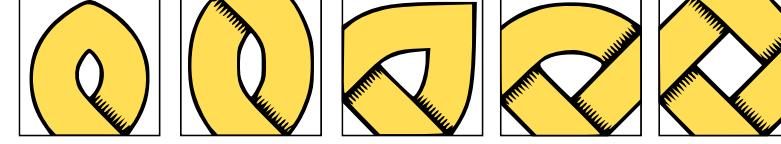
Board Pieces



n

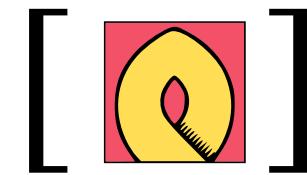
Problem: Can the player build a knot of size $\geq L$ given k pieces?

1-Player Celtic!: The “Easy” problems

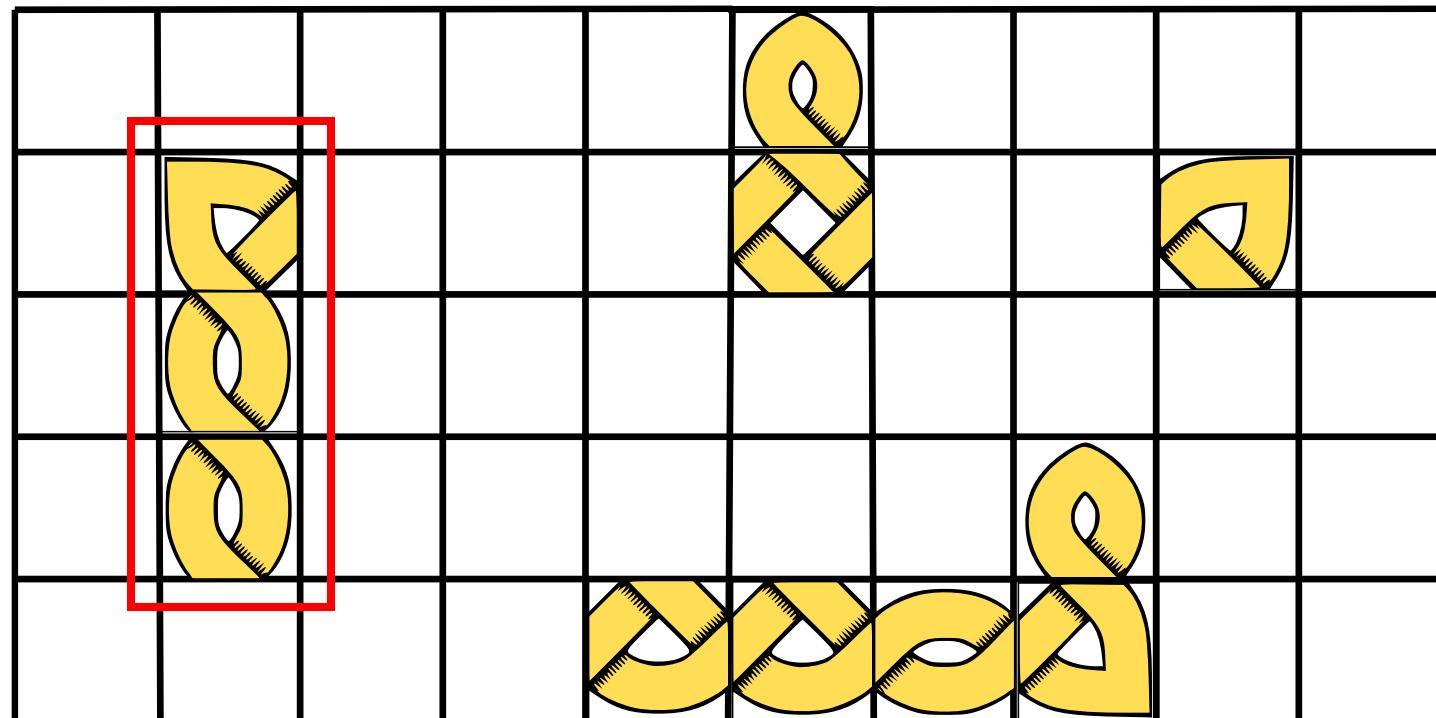
Player Pieces	Board Pieces	Complexity
		
		$O(n^2)$
		

Problem: Can the player build a knot of size $\geq L$ given k pieces?

1-Player Celtic!: The “Easy” problems

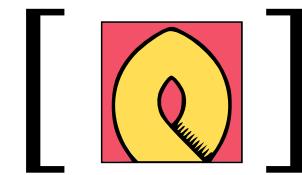


Player Pieces

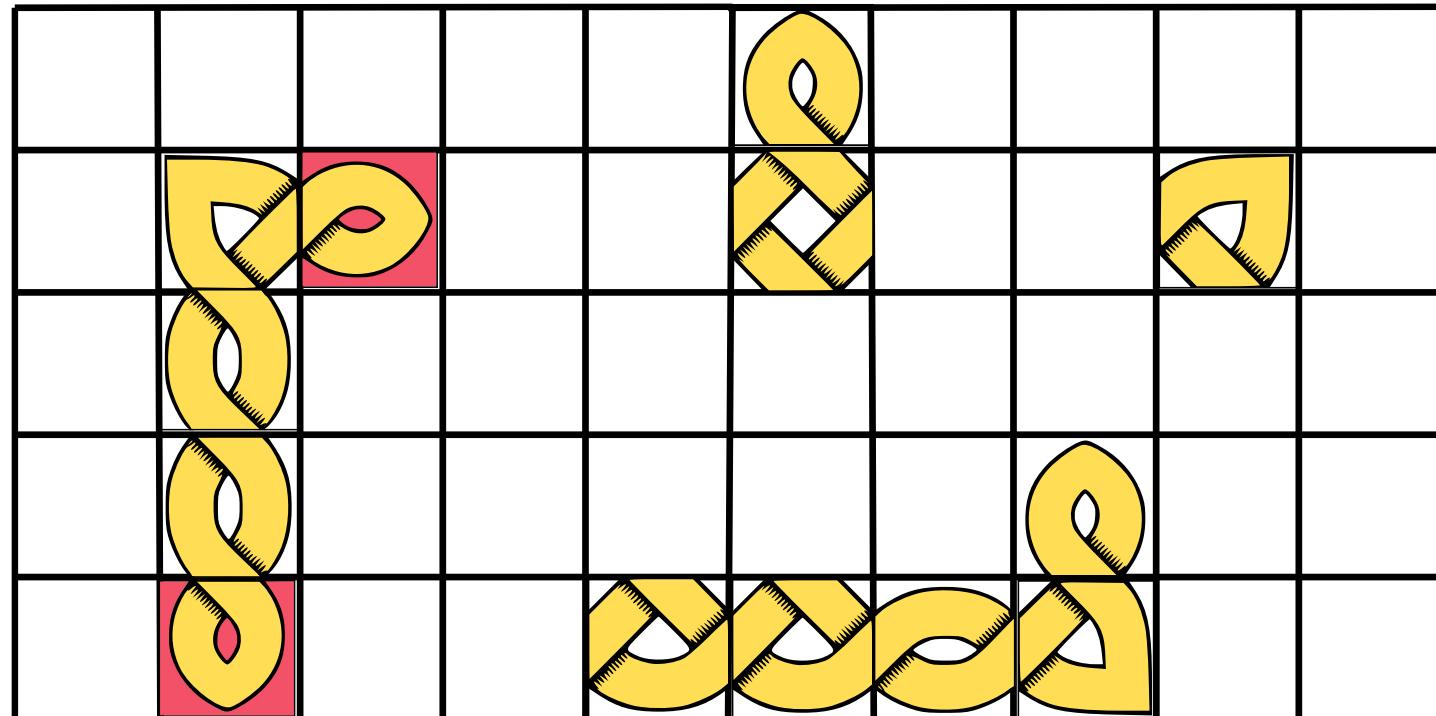


Given: $L = 6, k = 4$

1-Player Celtic!: The “Easy” problems

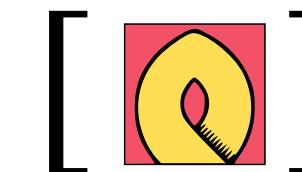


Player Pieces

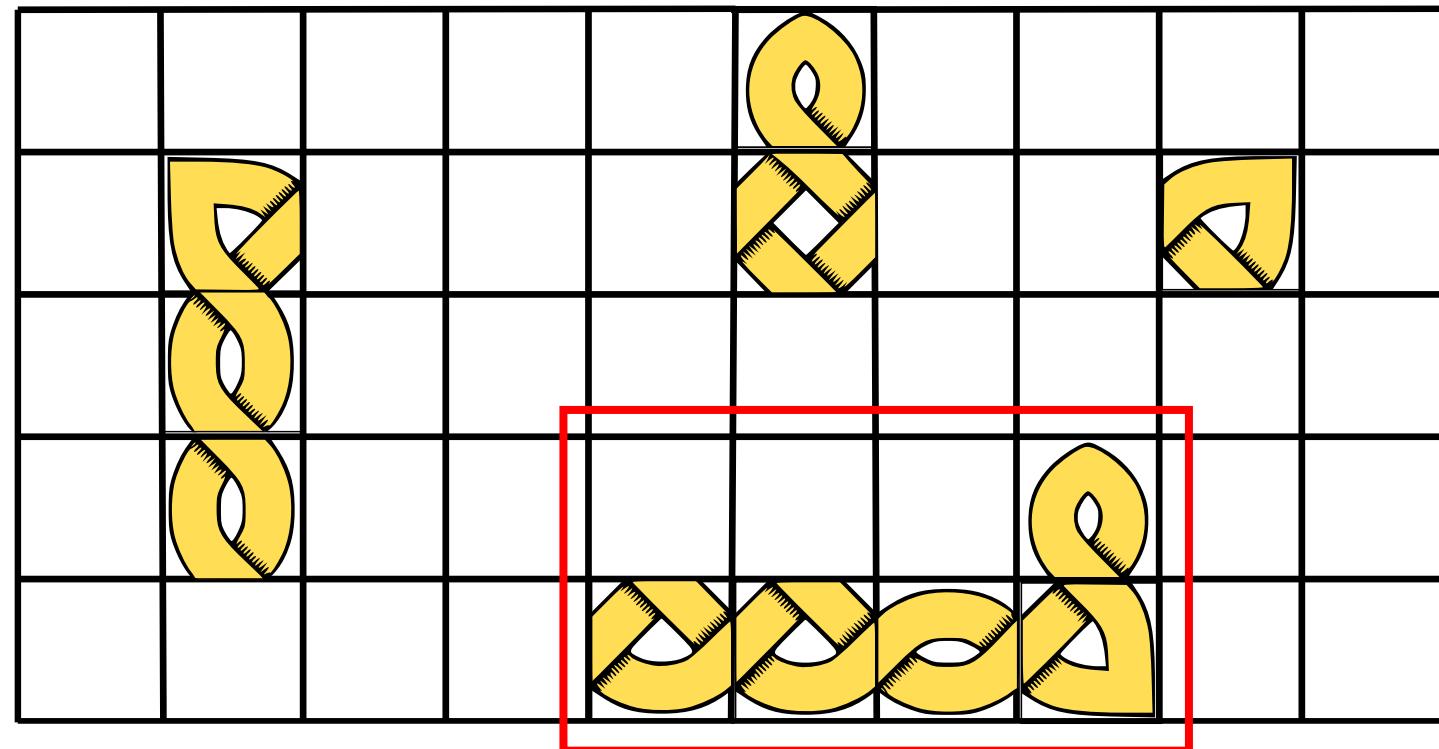


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1-Player Celtic!: The “Easy” problems

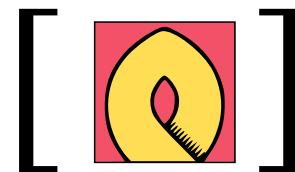


Player Pieces

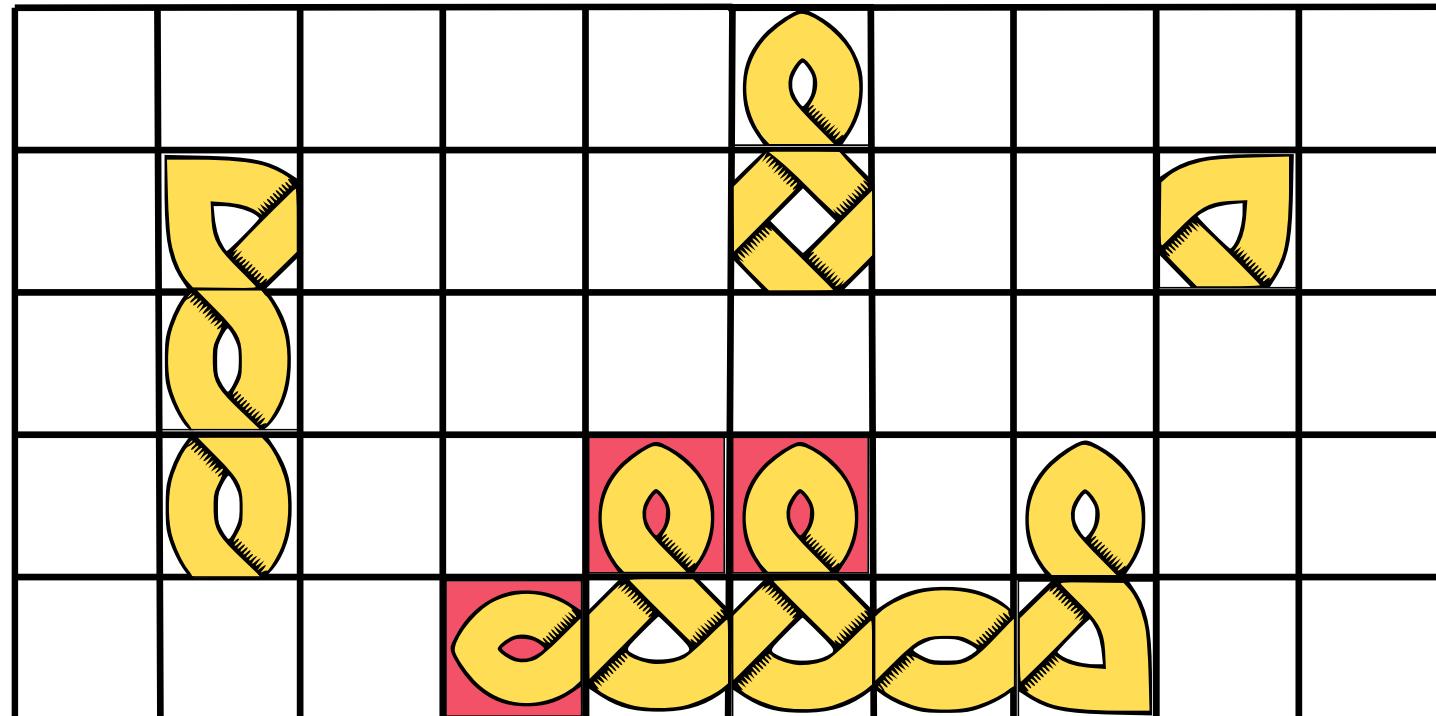


Given: $L = 6, k = 4$

1-Player Celtic!: The “Easy” problems

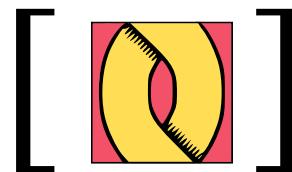


Player Pieces

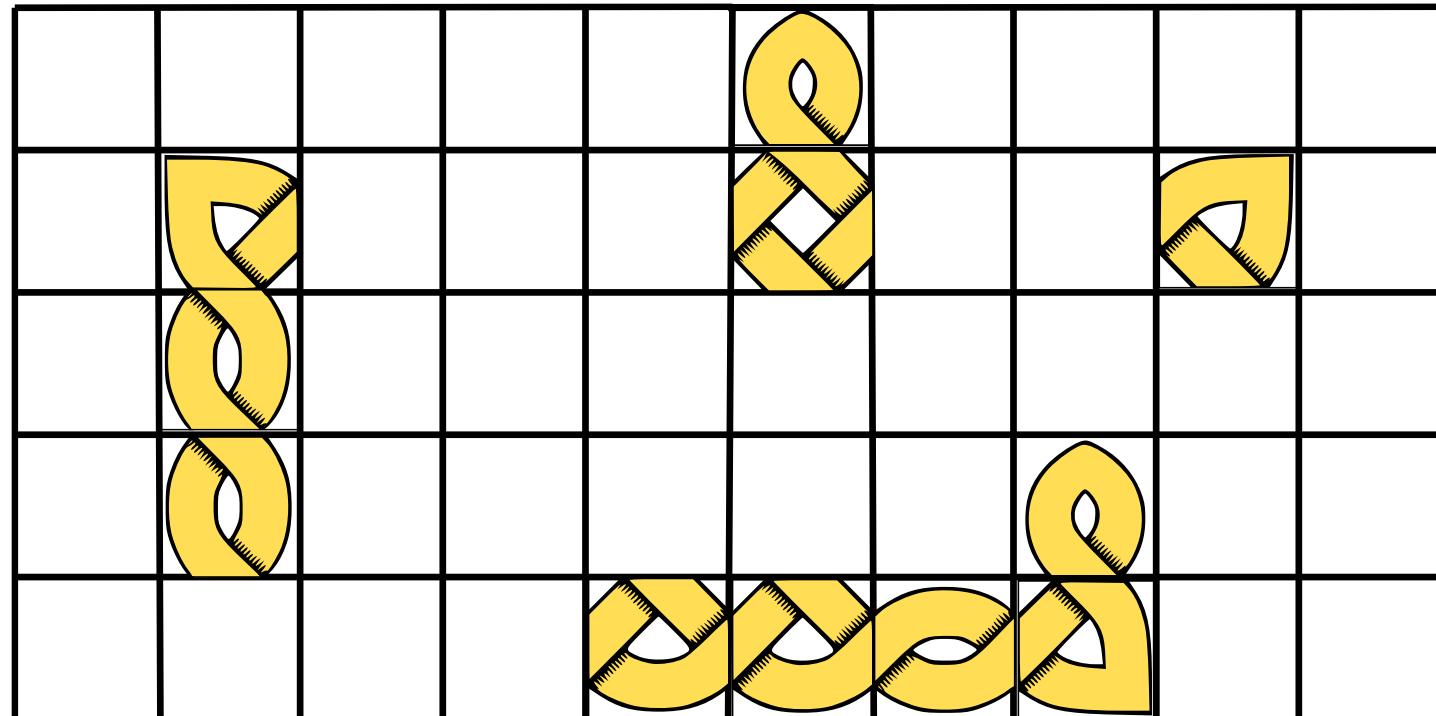


Given: $L = 6, k = 4$

1-Player Celtic!: The “Easy” problems

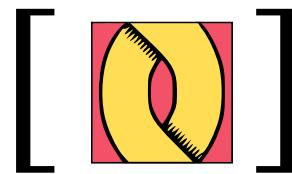


Player Pieces

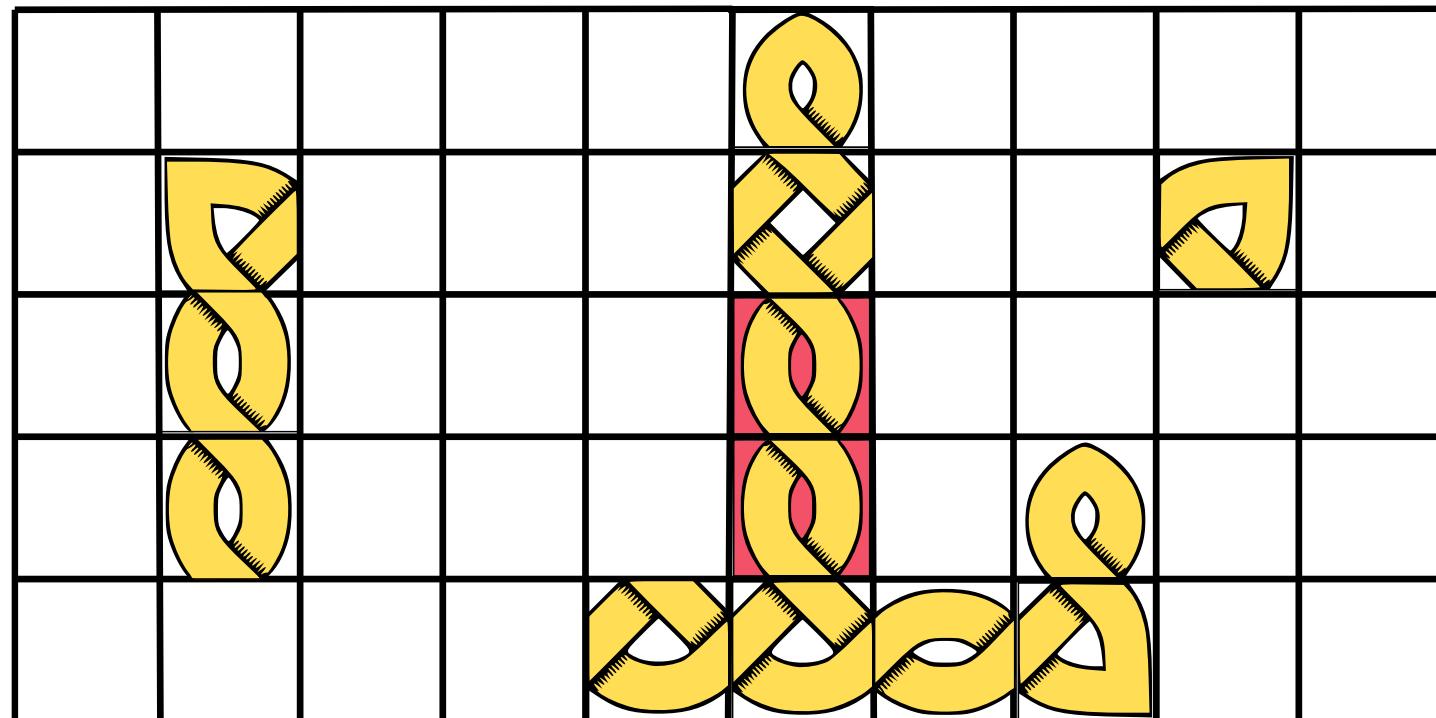


Given: $L = 6, k = 4$

1-Player Celtic!: The “Easy” problems

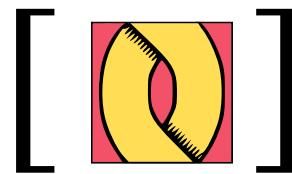


Player Pieces

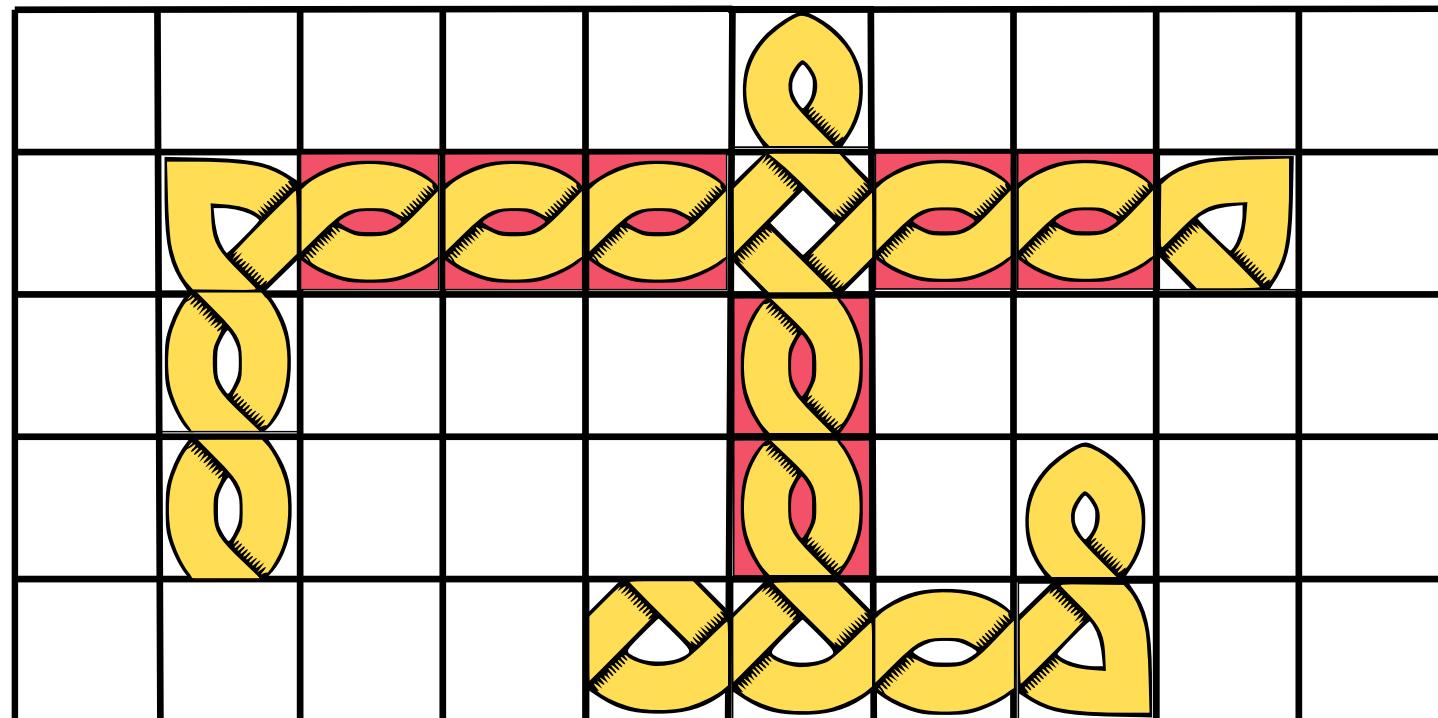


Given: $L = 6, k = 4$

1-Player Celtic!: The “Easy” problems

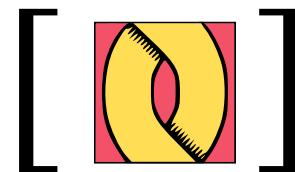


Player Pieces

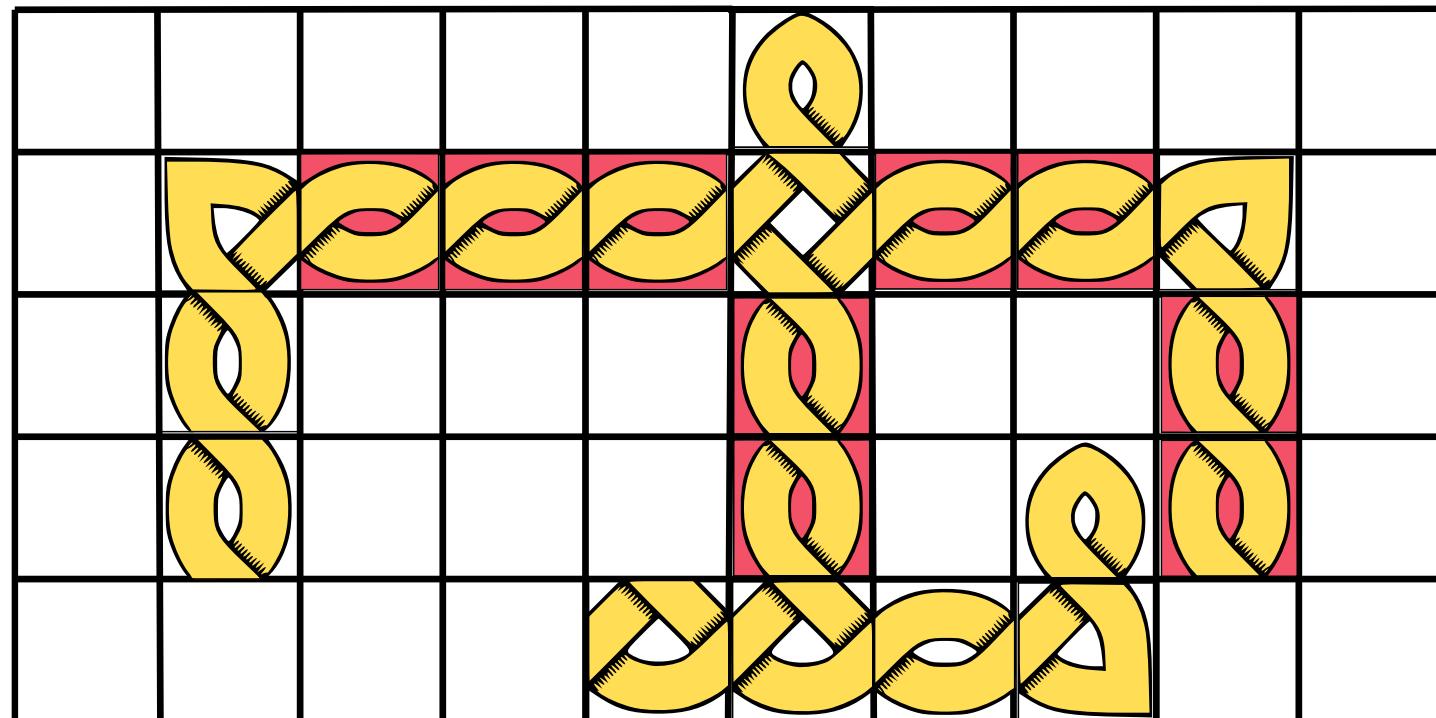


Given: $L = 6, k = 4$

1-Player Celtic!: The “Easy” problems



Player Pieces



Given: $L = 6, k = 4$

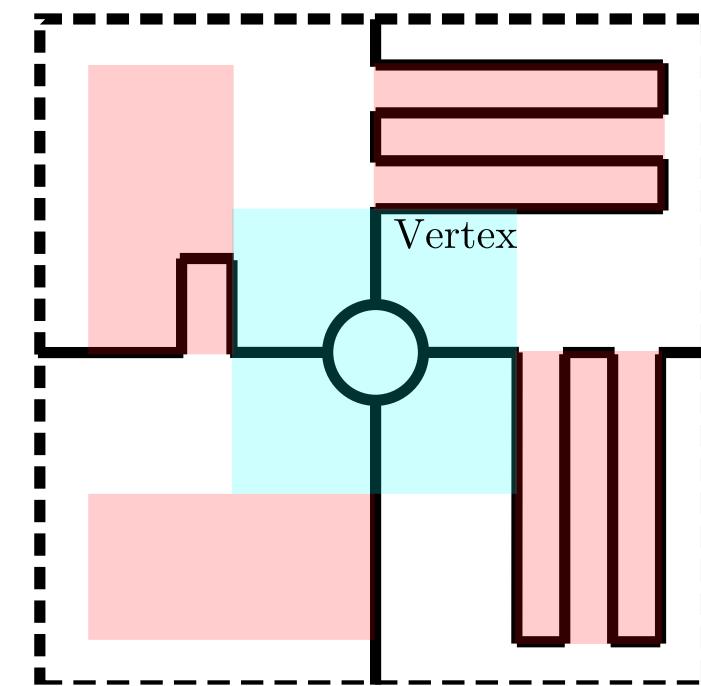
1-Player Celtic!: The “Hard” problems

Player Pieces	Board Pieces	Complexity
		<i>NP-complete</i>

Problem: Can the player build a knot of size $\geq L$ given k pieces?

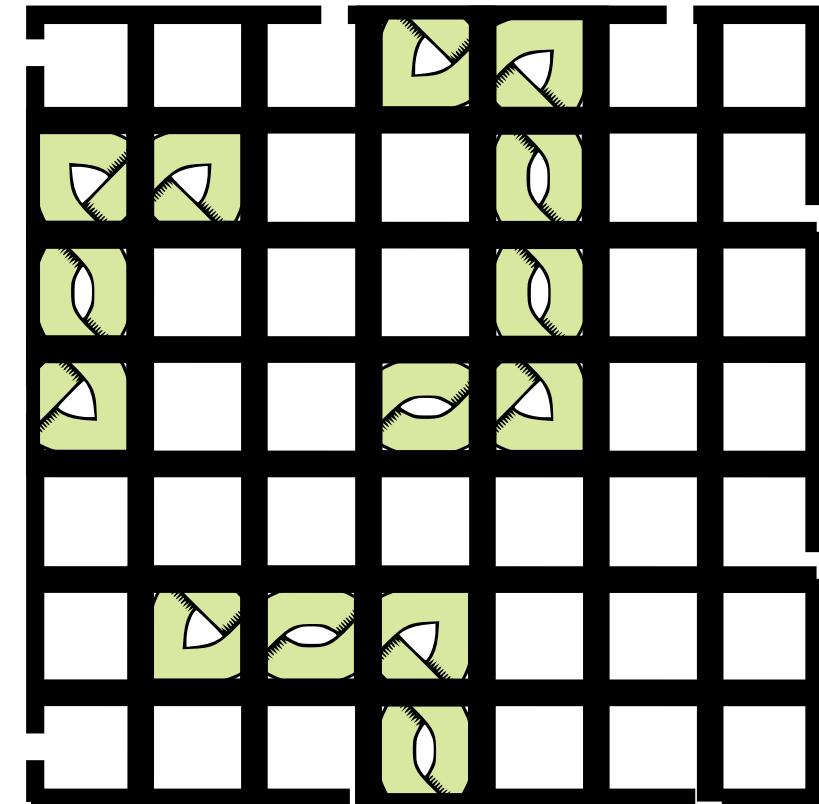
1-Player Celtic!: The “Hard” problems

- Planar, directed Hamiltonian Cycle with max degree 3
- Transform graph into equivalent rectilinear embedding with area $(|V| + 1) \times (|V| + 1)$
- Scale new graph by $O(|V|)$
- Edges are pumped to be roughly the same



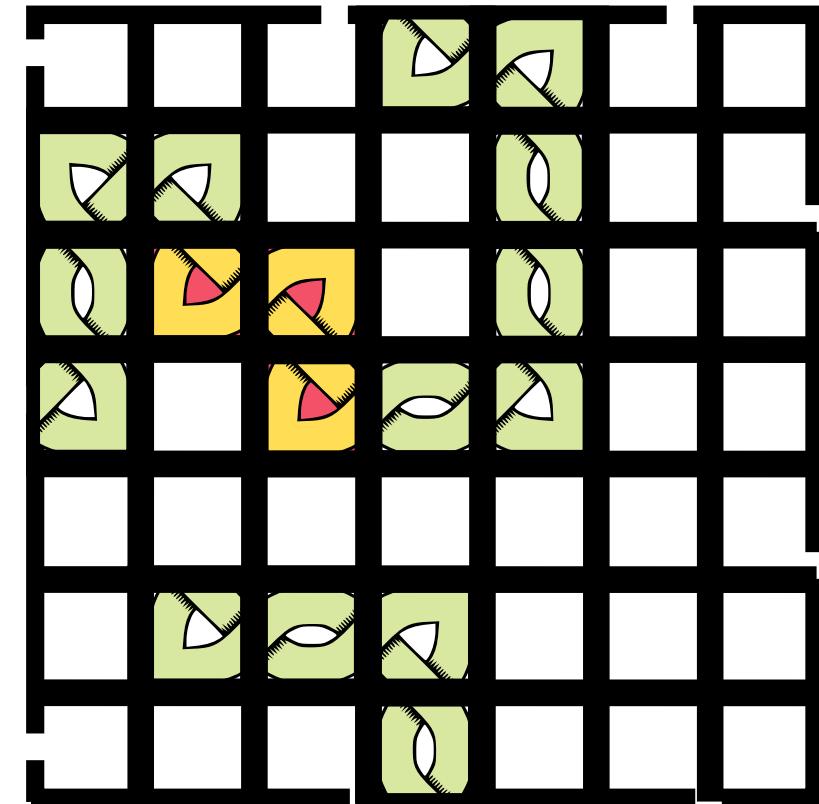
1-Player Celtic!: The “Hard” problems

- Center piece represents the “vertex”
- Incoming edge(s) are moved to one side
- Outgoing edge(s) are moved to the other side
- The “choice” is made by the played pieces

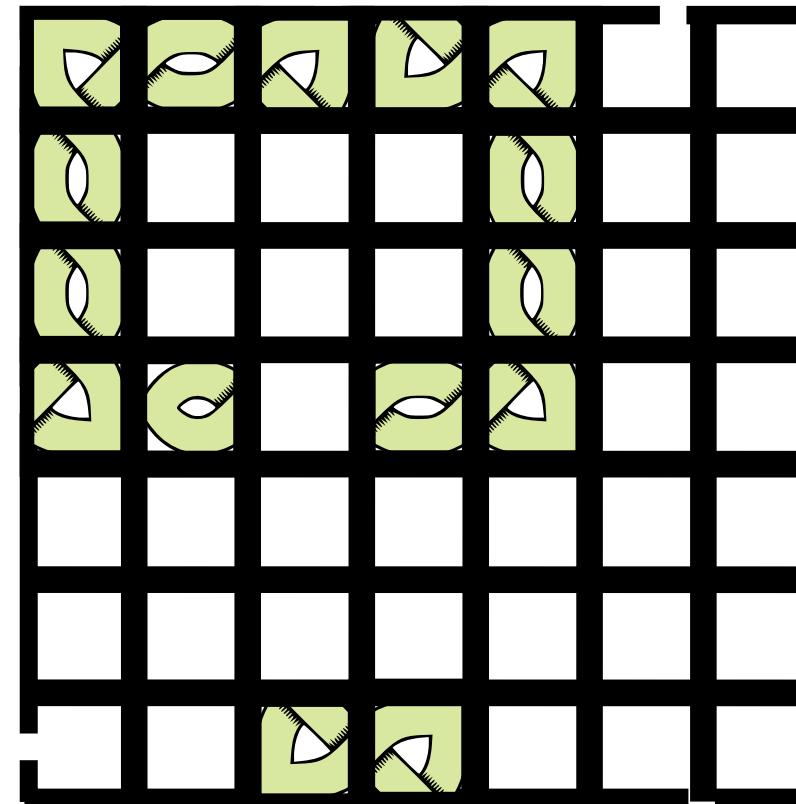
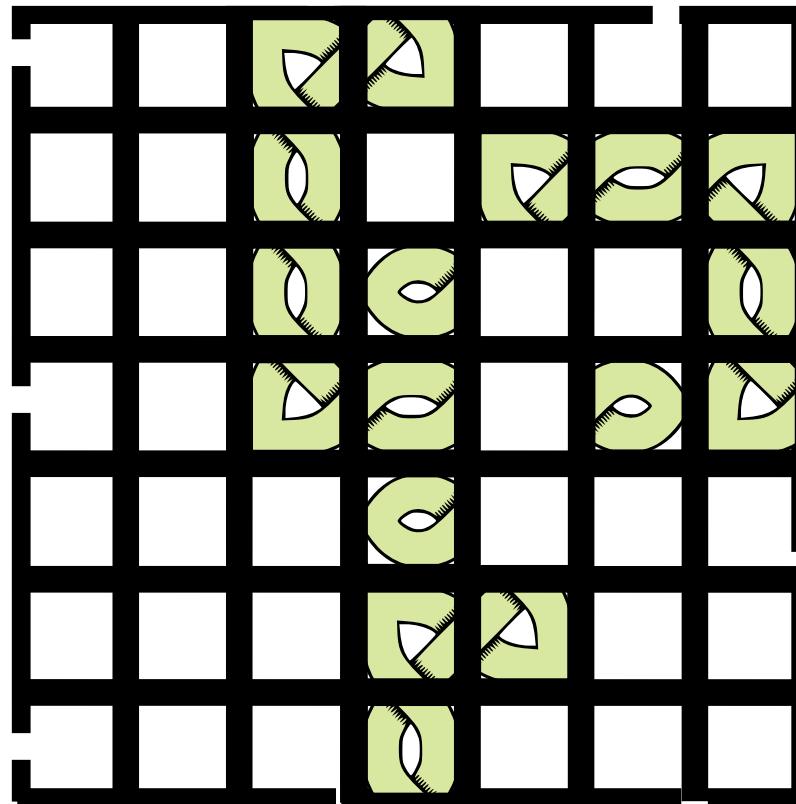


1-Player Celtic!: The “Hard” problems

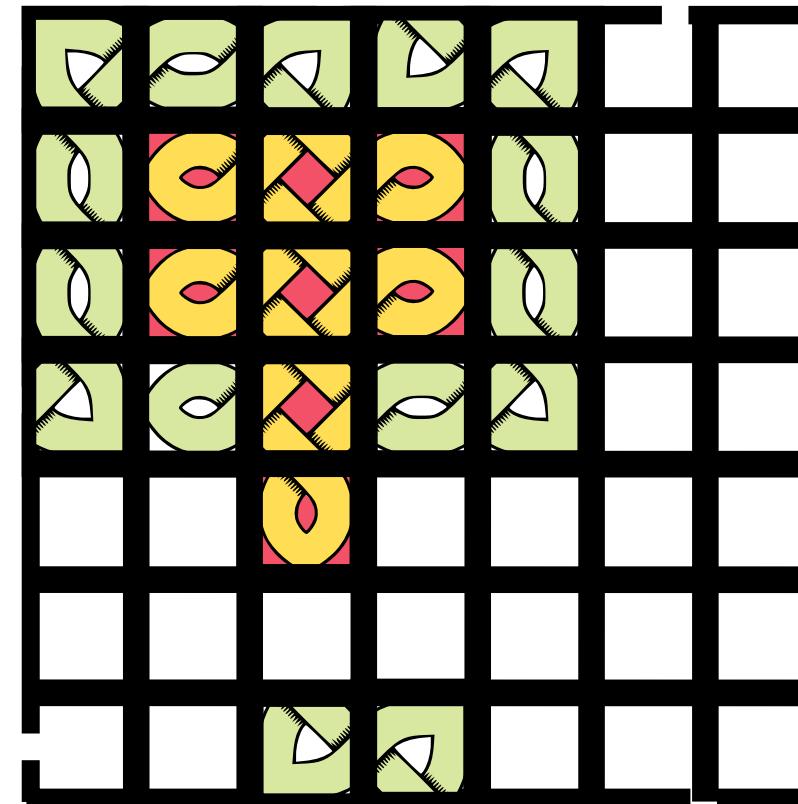
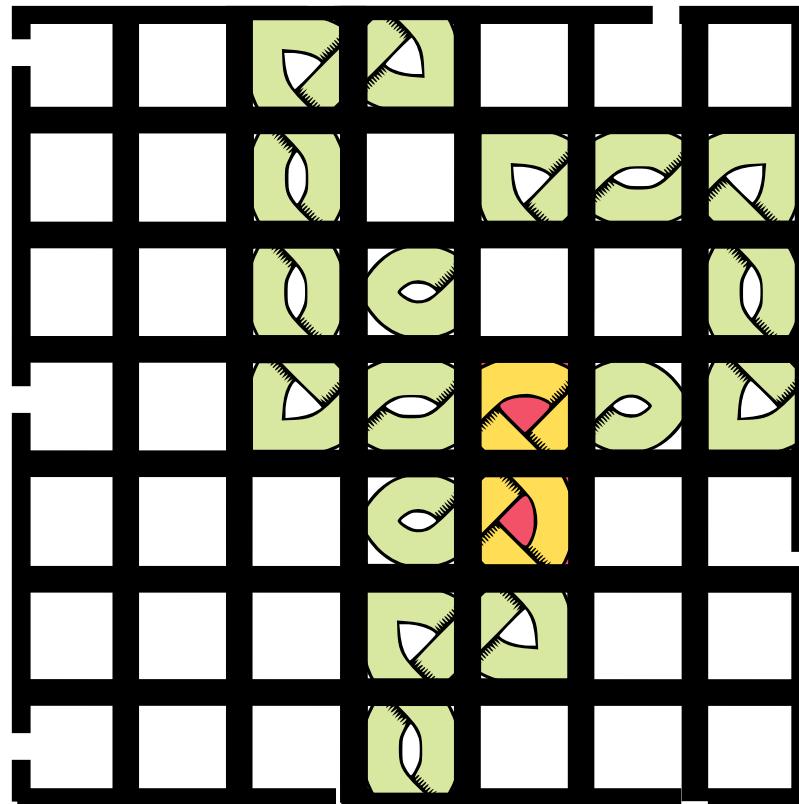
- Center piece represents the “vertex”
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1-Player Celtic!: The “Hard” problems

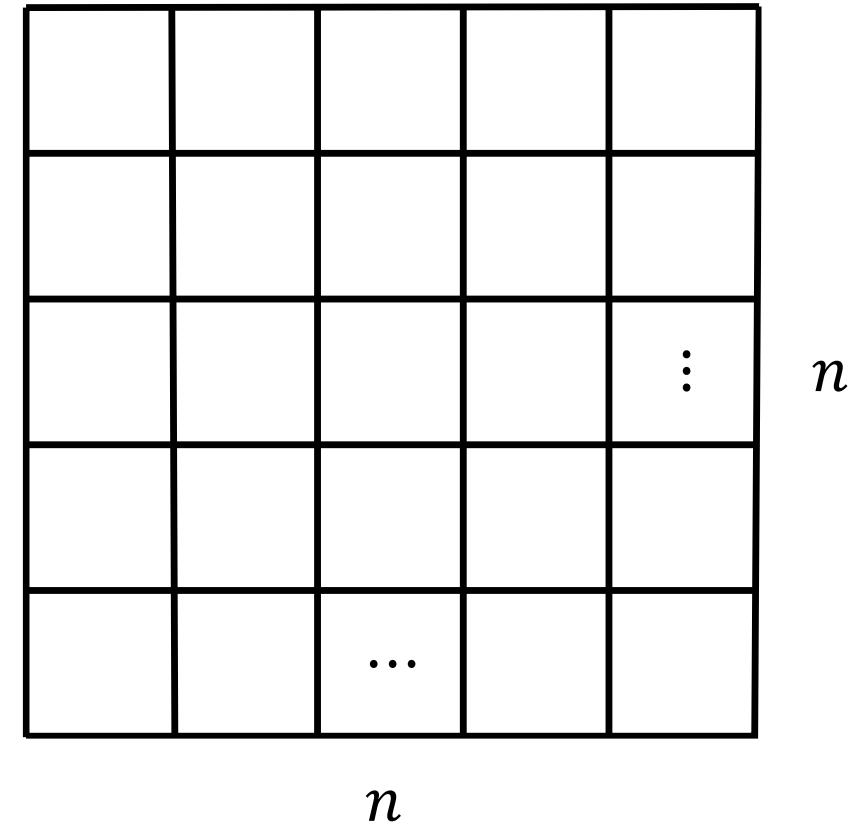
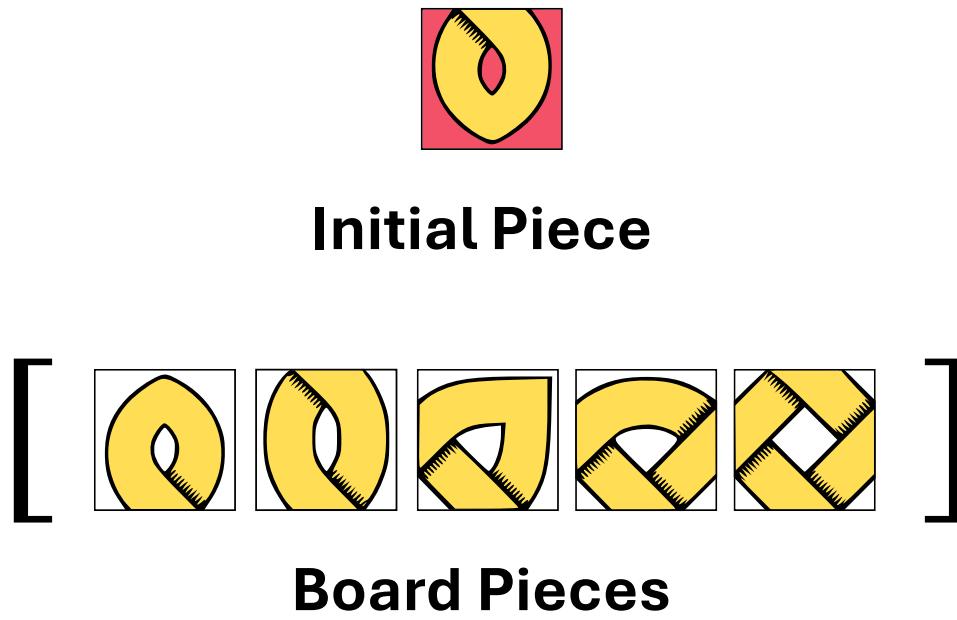


1-Player Celtic!: The “Hard” problems



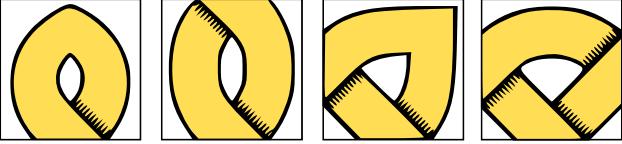
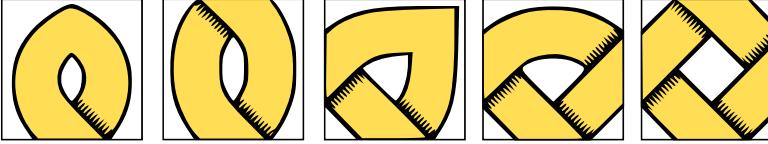
Generalized 0-Player Celtic!

Generalization: 0-Player



Problem: Does the simulation build a knot of size $\geq L$?

0-Player Celtic!

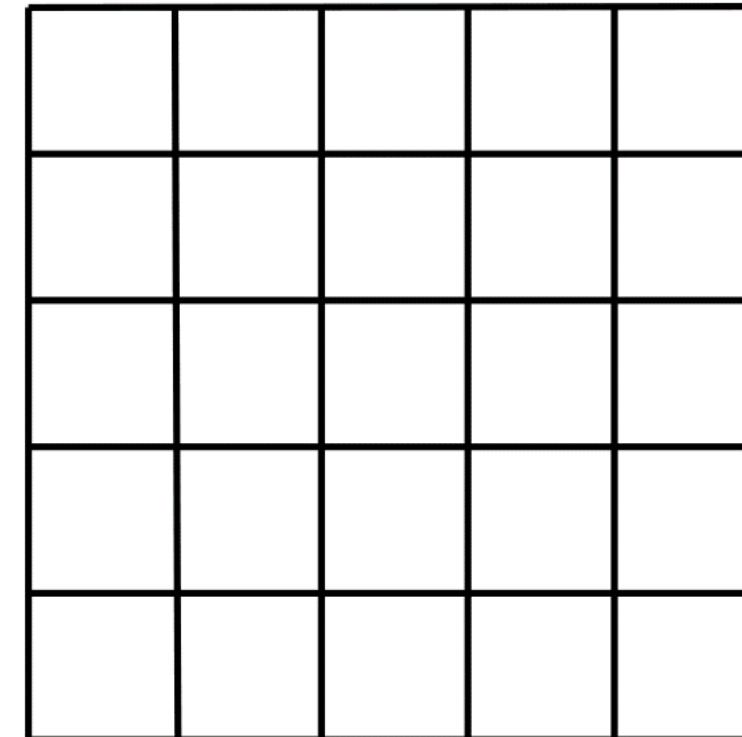
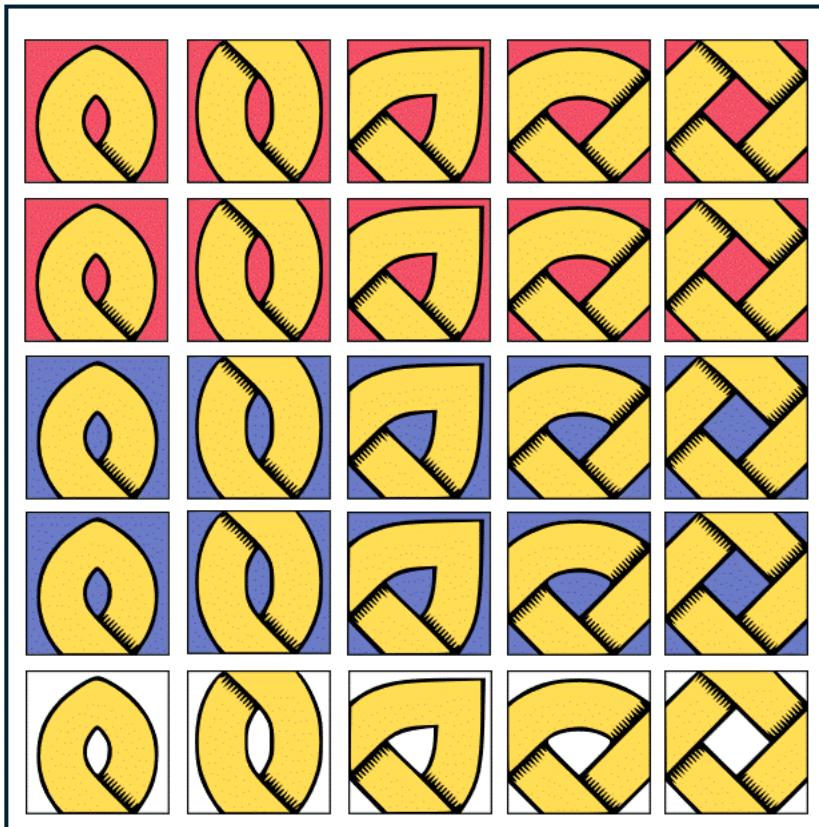
Player Pieces	Board Pieces	Complexity
		P -complete

Problem: Does the simulation build a knot of size $\geq L$?

Conclusion

- Some piece combinations missing for Generalized 1-Player Celtic!
- More balanced pieces for Generalized 2-Player Celtic!
- FPT, approximations for maximally sized knots
- Is Celtic! a 1st player win, 2nd player win, or a draw?
- Pattern based complexity of the knots.
- Consider > genus-1 knot assembly.

Tile-based Knot Assembly with Celtic!



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